

Sequent Systems and Provability Logic  $R^-$  for Rosser  
Sentences

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# Sequent Systems and Provability Logic $\mathbf{R}^-$ for Rosser Sentences

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## Abstract

To discuss Rosser sentences, Guaspari and Solovay [GS79] enriched the modal language by adding, for each  $\Box A$  and  $\Box B$ , the formulas  $\Box A \prec \Box B$  and  $\Box A \preceq \Box B$ , with arithmetic realizations. They introduced provability logics  $\mathbf{R}^-$ ,  $\mathbf{R}$  and  $\mathbf{R}^\omega$  with enriched language by extending the unimodal provability logic  $\mathbf{GL}$  and proved kinds of arithmetic completeness for them.

A sequent system for  $\mathbf{R}^-$ , the most preliminary logic among the logics they introduced, was given in Sasaki and Ohama [SO03]. They proved a cut-elimination theorem in weakened form, and as a result, a kind of subformula property was shown. However, considering a cut-free system for  $\mathbf{GL}$ , their system has a cut, which seems to be removable. Here we introduce another system with a kind of subformula property and discuss what kinds of cuts are removable from the system in [SO03]. Also we give a proof of completeness theorem without the extension lemma in [GS79], which was used in [SO03].

## 1 The logic $\mathbf{R}^-$

In this section, we introduce the logic  $\mathbf{R}^-$ . We use logical constant  $\perp$  (contradiction), and logical connectives  $\wedge$  (conjunction),  $\vee$  (disjunction),  $\supset$  (implication),  $\Box$  (provability),  $\preceq$  (witness comparison), and  $\prec$  (witness comparison). Formulas are defined inductively as follows:

- (1) propositional variables and  $\perp$  are formulas,
  - (2) if  $A$  and  $B$  are formulas, then so are  $(A \wedge B)$ ,  $(A \vee B)$ ,  $(A \supset B)$ ,  $(\Box A)$ ,  $(\Box A \prec \Box B)$  and  $(\Box A \preceq \Box B)$ .
- A formula of the form  $\Box A$  is said to be a  $\Box$ -formula. Also a formula of the form  $\Box A \preceq \Box B$  ( $\Box A \prec \Box B$ ) is said to be a  $\preceq$ -formula ( $\prec$ -formula). By  $\Sigma$ , we mean the set of all  $\Box$ -formulas, all  $\prec$ -formulas and all  $\preceq$ -formulas.

The modal system  $\mathbf{R}^-$  is defined by the following axioms and inference rules.

### Axioms of $\mathbf{R}^-$

- A1 : all tautologies,
- A2 :  $\Box(A \supset B) \supset (\Box A \supset \Box B)$ ,
- A3 :  $\Box(\Box A \supset A) \supset \Box A$ ,
- A4 :  $A \supset \Box A$ , where  $A \in \Sigma$ ,
- A5 :  $(\Box A \preceq \Box B) \supset \Box A$ ,
- A6 :  $((\Box A \preceq \Box B) \wedge (\Box B \preceq \Box C)) \supset (\Box A \preceq \Box C)$ ,
- A7 :  $(\Box A \vee \Box B) \supset ((\Box A \preceq \Box B) \vee (\Box B \prec \Box A))$ ,
- A8 :  $(\Box A \prec \Box B) \supset (\Box A \preceq \Box B)$ ,
- A9 :  $((\Box A \preceq \Box B) \wedge (\Box B \prec \Box A)) \supset \perp$ ,

### Inference rules of $\mathbf{R}^-$

- MP :  $A, A \supset B \in \mathbf{R}^-$  implies  $B \in \mathbf{R}^-$ ,
- N :  $A \in \mathbf{R}^-$  implies  $\Box A \in \mathbf{R}^-$ .

In [GS79] and Smorínski [Smo85], the following two formulas are also axioms of  $\mathbf{R}^-$ , but they are redundant (cf. De Jongh [Jon87] and Voorbraak [Voo90]).

- A10 :  $\Box A \supset (\Box A \preceq \Box A)$ ,
- A11 :  $(\Box A \wedge (\Box B \supset \perp)) \supset (\Box A \prec \Box B)$ .

We introduce Kripke semantics for  $\mathbf{R}^-$ , following [Smo85].<sup>1</sup>

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<sup>1</sup>[Smo85] uses a rooted frame, but our frame does not necessarily be rooted.

**Definition 1.1.** A Kripke pseudo-model for  $\mathbf{R}^-$  is a triple  $\langle \mathbf{W}, <, \models \rangle$  where

- (1)  $\mathbf{W}$  is a non-empty finite set,
- (2)  $<$  is an irreflexive and transitive binary relation on  $\mathbf{W}$  satisfying

$$\alpha < \gamma \text{ and } \beta < \gamma \text{ imply either one of } \alpha = \beta, \alpha < \beta \text{ or } \beta < \alpha,$$

- (3)  $\models$  is a valuation satisfying, in addition to the usual boolean laws,

$$\alpha \models \Box A \text{ if and only if for any } \beta \in \alpha \uparrow (= \{\gamma \mid \alpha < \gamma\}), \beta \models A.$$

**Definition 1.2.** A Kripke pseudo-model  $\langle \mathbf{W}, <, \models \rangle$  for  $\mathbf{R}^-$  is said to be a Kripke model for  $\mathbf{R}^-$  if the following conditions hold, for any formula  $D$ ,

- (1) if  $D \in \Sigma$  and  $\alpha \models D$ , then for any  $\beta \in \alpha \uparrow, \beta \models D$ ,
- (2) if  $D$  is either one of the axioms  $A5, A6, A7, A8$  and  $A9$ , then  $\alpha \models D$ .

**Lemma 1.3**([GS79]).  $A \in \mathbf{R}^-$  if and only if  $A$  is valid in any Kripke model for  $\mathbf{R}^-$ .

## 2 A sequent system $\mathbf{GR}^-$

In this section we introduce a sequent system  $\mathbf{GR}^-$  for  $\mathbf{R}^-$ . We use Greek letters,  $\Gamma$  and  $\Delta$ , possibly with suffixes, for finite sets of formulas. The expression  $\Box \Gamma$  denotes the set  $\{\Box A \mid A \in \Gamma\}$ . By a sequent, we mean the expression  $\Gamma \rightarrow \Delta$ . For brevity's sake, we write

$$A_1, \dots, A_k, \Gamma_1, \dots, \Gamma_\ell \rightarrow \Delta_1, \dots, \Delta_m, B_1, \dots, B_n$$

instead of

$$\{A_1, \dots, A_k\} \cup \Gamma_1 \cup \dots \cup \Gamma_\ell \rightarrow \Delta_1 \cup \dots \cup \Delta_m \cup \{B_1, \dots, B_n\}.$$

By  $\text{Sub}(A)$ , we mean the set of subformulas of  $A$ . We put

$$\text{Sub}^+(A) = \text{Sub}(A) \cup \{\Box B \preceq \Box C \mid \Box B, \Box C \in \text{Sub}(A)\} \cup \{\Box B \prec \Box C \mid \Box B, \Box C \in \text{Sub}(A)\},$$

$$\text{Sub}(\Gamma \rightarrow \Delta) = \bigcup_{B \in \Gamma \cup \Delta} \text{Sub}(B), \quad \text{Sub}^+(\Gamma \rightarrow \Delta) = \bigcup_{B \in \Gamma \cup \Delta} \text{Sub}^+(B).$$

By the sequent system  $\mathbf{LK}$  for the classical propositional logic, we mean the system defined by the following axioms and inference rules in the usual way.

**Axioms of  $\mathbf{LK}$ :**

$$A \rightarrow A$$

$$\perp \rightarrow$$

**Inference rules of  $\mathbf{LK}$ :**

$$\frac{\Gamma \rightarrow \Delta}{A, \Gamma \rightarrow \Delta} (w \rightarrow) \qquad \frac{\Gamma \rightarrow \Delta}{\Gamma \rightarrow \Delta, A} (\rightarrow w)$$

$$\frac{\Gamma \rightarrow \Delta, A \quad A, \Pi \rightarrow \Lambda}{\Gamma, \Pi \rightarrow \Delta, \Lambda} (cut)$$

$$\frac{A_i, \Gamma \rightarrow \Delta}{A_1 \wedge A_2, \Gamma \rightarrow \Delta} (\wedge \rightarrow_i) \qquad \frac{\Gamma \rightarrow \Delta, A \quad \Gamma \rightarrow \Delta, B}{\Gamma \rightarrow \Delta, A \wedge B} (\rightarrow \wedge)$$

$$\frac{A, \Gamma \rightarrow \Delta \quad B, \Gamma \rightarrow \Delta}{A \vee B, \Gamma \rightarrow \Delta} (\vee \rightarrow) \qquad \frac{\Gamma \rightarrow \Delta, A_i}{\Gamma \rightarrow \Delta, A_1 \vee A_2} (\rightarrow \vee_i)$$

$$\frac{\Gamma \rightarrow \Delta, A \quad B, \Gamma \rightarrow \Delta}{A \supset B, \Gamma \rightarrow \Delta} (\supset \rightarrow) \qquad \frac{A, \Gamma \rightarrow \Delta, B}{\Gamma \rightarrow \Delta, A \supset B} (\rightarrow \supset)$$

The system  $\mathbf{GR}^-$  is obtained from the sequent system  $\mathbf{LK}$  by adding the following axioms and inference rules in the usual way.

**Additional axioms of  $\mathbf{GR}^-$**

- GA1:  $\Box A \preceq \Box B, \Box B \preceq \Box C \rightarrow \Box A \preceq \Box C$
- GA2:  $\Box A \rightarrow \Box A \preceq \Box B, \Box B \prec \Box A$
- GA3:  $\Box B \rightarrow \Box A \preceq \Box B, \Box B \prec \Box A$
- GA4:  $\Box A \prec \Box B \rightarrow \Box A \preceq \Box B$
- GA5:  $\Box A \preceq \Box B, \Box B \prec \Box A \rightarrow$

**Additional inference rules of  $\mathbf{GR}^-$**

$$\frac{\Box A, \Sigma^f, \Gamma \rightarrow A}{\Sigma^f, \Box \Gamma \rightarrow \Box A} (\rightarrow \Box) \quad \frac{\Box A, \Gamma \rightarrow \Delta}{\Box A \preceq \Box B, \Gamma \rightarrow \Delta} (\preceq \rightarrow) \quad \frac{\Gamma \rightarrow \Delta, \Box A}{\Gamma \rightarrow \Delta, \Box A \preceq \Box A} (\rightarrow \preceq)$$

where  $\Sigma^f$  is a finite subset of  $\Sigma$ .

The system  $\mathbf{GR}_1^-$  is the system obtained from  $\mathbf{GR}^-$  by restricting a cut to the following form:

$$\frac{\Gamma \rightarrow \Delta, \Box A \odot \Box B \quad \Box A \odot \Box B, \Gamma \rightarrow \Delta}{\Gamma \rightarrow \Delta}$$

where  $\odot \in \{\prec, \preceq\}$ , and  $\Box A$  and  $\Box B$  are subformulas of a formula occurring in the lower sequent.

**Theorem 2.1**([SO03]). *The following conditions are equivalent:*

- (1)  $A_1, \dots, A_m \rightarrow B_1, \dots, B_n \in \mathbf{GR}_1^-$ ,
- (2)  $A_1, \dots, A_m \rightarrow B_1, \dots, B_n \in \mathbf{GR}^-$ ,
- (3)  $A_1 \wedge \dots \wedge A_m \supset B_1 \vee \dots \vee B_n \in \mathbf{R}^-$ ,
- (4)  $A_1 \wedge \dots \wedge A_m \supset B_1 \vee \dots \vee B_n$  is valid in any Kripke model for  $\mathbf{R}^-$ .

**Corollary 2.2.** *If a sequent  $S$  is provable in  $\mathbf{GR}^-$ , then there exists a proof figure  $\mathcal{P}$  for  $S$  such that each formula occurring in  $\mathcal{P}$  belongs to  $\text{Sub}^+(S)$ .*

### 3 A sequent system for $\mathbf{R}^-$ without additional axioms

Theorem 2.1 provides the decision procedure for the provability of  $\mathbf{R}^-$ , but does not say that every cut in  $\mathbf{GR}_1^-$  is necessary. For instance, the following cut seems to be removable if  $\Gamma \rightarrow \Delta$  does not have any  $\prec$ -formula and  $\preceq$ -formula. Because we can easily see the  $\prec$ -free and  $\preceq$ -free fragment of  $\mathbf{GR}^-$  is the system for the provability logic  $\mathbf{GL}$ , the  $\prec$ -free and  $\preceq$ -free fragment of  $\mathbf{R}^-$ , described in Valentini [Val83] and Avron [Avr84] and enjoying a cut-elimination theorem.

$$\frac{\frac{\Gamma \rightarrow \Delta, \Box A}{\Gamma \rightarrow \Delta, \Box A \preceq \Box A} (\rightarrow \preceq) \quad \frac{\Box A, \Gamma \rightarrow \Delta}{\Box A \preceq \Box A, \Gamma \rightarrow \Delta} (\preceq \rightarrow)}{\Gamma \rightarrow \Delta} (cut)$$

Here we introduce another system for  $\mathbf{R}^-$  by adding only inference rules to  $\mathbf{LK}$  and prove cut-elimination theorem. Also we consider what kind of cuts are removable from  $\mathbf{GR}_1^-$ .

The system  $\mathbf{GR}_2^-$  is obtained from  $\mathbf{LK}$  by adding the following inference rules in the usual way.

**Additional inference rules of  $\mathbf{GR}_2^-$**

$$(\rightarrow \Box), (\rightarrow \preceq), (\preceq \rightarrow) \text{ are as in } \mathbf{GR}^-,$$

$$\begin{array}{c}
\frac{\Box A \preceq \Box B, \Gamma \rightarrow \Delta, \Box B \preceq \Box A}{\Box A \prec \Box B, \Gamma \rightarrow \Delta} (\prec \rightarrow) \\
\\
\frac{\Gamma \rightarrow \Delta, \Box C \preceq \Box D \quad \Gamma \rightarrow \Delta, \Box D \preceq \Box E \quad \Box C \preceq \Box E, \Gamma \rightarrow \Delta}{\Gamma \rightarrow \Delta} (tran) \\
\\
\frac{\Gamma \rightarrow \Delta, \Box C, \Box D \quad \Box C \prec \Box D, \Gamma \rightarrow \Delta \quad \Box D \prec \Box C, \Gamma \rightarrow \Delta \quad \Box C \preceq \Box D, \Box D \preceq \Box C, \Gamma \rightarrow \Delta}{\Gamma \rightarrow \Delta} (lin)
\end{array}$$

where  $\Box C, \Box D$  and  $\Box E$  are different subformulas occurring in the lower sequent.

The system  $\mathbf{GR}_3^-$  is the system obtained from  $\mathbf{GR}_2^-$  by removing cuts.

**Lemma 3.1.** *The following conditions are equivalent:*

- (1)  $\Gamma \rightarrow \Delta \in \mathbf{GR}^-$ ,
- (2)  $\Gamma \rightarrow \Delta \in \mathbf{GR}_2^-$ .

Proof. For “(1) implies (2)”. Additional inference rules of  $\mathbf{GR}^-$  are also inference rule in  $\mathbf{GR}_2^-$ . So, it is sufficient to show the provability of the additional axioms of  $\mathbf{GR}^-$  in  $\mathbf{GR}_2^-$ . Axioms  $GA4(\Box A \prec \Box B \rightarrow \Box A \preceq \Box B)$  and  $GA5(\Box A \preceq \Box B, \Box B \prec \Box A \rightarrow)$  are shown by the following figures.

$$\begin{array}{c}
\frac{\Box A \preceq \Box B \rightarrow \Box A \preceq \Box B}{\Box A \preceq \Box B \rightarrow \Box A \preceq \Box B, \Box B \prec \Box A} (\rightarrow w) \\
\frac{\Box A \preceq \Box B \rightarrow \Box A \preceq \Box B, \Box B \prec \Box A}{\Box A \prec \Box B \rightarrow \Box A \preceq \Box B} (\prec \rightarrow) \\
\\
\frac{\Box A \preceq \Box B \rightarrow \Box A \preceq \Box B}{\Box B \preceq \Box A, \Box A \preceq \Box B \rightarrow \Box A \preceq \Box B} (\rightarrow w) \\
\frac{\Box B \preceq \Box A, \Box A \preceq \Box B \rightarrow \Box A \preceq \Box B}{\Box A \preceq \Box B, \Box B \prec \Box A \rightarrow} (\prec \rightarrow)
\end{array}$$

For  $GA1(\Box A \preceq \Box B, \Box B \preceq \Box C \rightarrow \Box A \preceq \Box C)$ . If  $A, B$  and  $C$  are different, then the provability can be shown by *(tran)* and weakening rules. If  $A = B$  or  $B = C$ , then it can be shown by weakening rules. If  $A = C$ , by the following figure.

$$\begin{array}{c}
\frac{\Box A \rightarrow \Box A}{\Box A \rightarrow \Box A \preceq \Box A} (\rightarrow \preceq) \\
\frac{\Box A \rightarrow \Box A \preceq \Box A}{\Box A \preceq \Box B \rightarrow \Box A \preceq \Box A} (\preceq \rightarrow) \\
\frac{\Box A \preceq \Box B \rightarrow \Box A \preceq \Box A}{\Box A \preceq \Box B, \Box B \preceq \Box A \rightarrow \Box A \preceq \Box A} (w \rightarrow)
\end{array}$$

For  $GA2(\Box A \rightarrow \Box A \preceq \Box B, \Box B \prec \Box A)$ . If  $A \neq B$ , then the provability can be shown by *(lin)*, the provability of  $GA4$  and weakening rules. If  $A = B$ , then it can be shown by  $(\rightarrow \preceq)$  and weakening rules. The provability of  $GA3$  can be shown similarly to  $GA2$ .

For “(2) implies (1)”. By the figures in the next page, each inference rule in  $\mathbf{GR}_2^-$  preserves the provability of  $\mathbf{GR}^-$ . ⊥

**Theorem 3.2.** *The following conditions are equivalent:*

- (1)  $A_1, \dots, A_m \rightarrow B_1, \dots, B_n \in \mathbf{GR}_3^-$ ,
- (2)  $A_1, \dots, A_m \rightarrow B_1, \dots, B_n \in \mathbf{GR}_2^-$ .

“(1) implies (2)” is clear. To prove “(2) implies (1)”, we need some preparations.

$(\prec \rightarrow)$ :

$$\frac{\frac{\frac{\Box A \prec \Box B, \Gamma \rightarrow \Delta, \Box B \preceq \Box A}{\Box A \prec \Box B, \Box A \preceq \Box B} \quad \frac{\Box A \preceq \Box B, \Box A \prec \Box B, \Gamma \rightarrow \Delta, \Box B \preceq \Box A}{\Box A \prec \Box B, \Gamma \rightarrow \Delta, \Box B \preceq \Box A} \quad (w \rightarrow)}{\Box B \preceq \Box A, \Box A \prec \Box B \rightarrow} \quad (cut)$$

$(tran)$ :

$$\frac{\frac{\Gamma \rightarrow \Delta, \Box D \preceq \Box E \quad \Box D \preceq \Box E, \Box C \preceq \Box D \rightarrow \Box C \preceq \Box E}{\Box C \preceq \Box D, \Gamma \rightarrow \Delta, \Box C \preceq \Box E} \quad (cut)}{\Gamma \rightarrow \Delta, \Box C \preceq \Box E} \quad (cut) \quad \frac{\Box C \preceq \Box E, \Gamma \rightarrow \Delta}{\Gamma \rightarrow \Delta} \quad (cut)$$

$(lim)$ :

$$\frac{\frac{\frac{\Gamma \rightarrow \Delta, \Box C, \Box D}{\Gamma \rightarrow \Delta, \Box C, \Box D \preceq \Box D} \quad (\rightarrow \preceq) \quad \mathcal{P}(D, C)}{\Gamma \rightarrow \Delta, \Box C \preceq \Box D, \Box D \preceq \Box C, \Box C} \quad (cut)}{\Gamma \rightarrow \Delta, \Box C \preceq \Box D, \Box D \preceq \Box C, \Box C \preceq \Box C} \quad (\rightarrow \preceq) \quad \mathcal{P}(C, D)} \quad (cut) \quad \frac{\mathcal{Q}(C, D)}{\Gamma \rightarrow \Delta, \Box C \preceq \Box D, \Box D \preceq \Box C} \quad (cut) \quad \frac{\mathcal{Q}(D, C) \quad \Box D \preceq \Box C, \Box C \preceq \Box D, \Gamma \rightarrow \Delta}{\Box C \preceq \Box D, \Gamma \rightarrow \Delta} \quad (cut)}{\Gamma \rightarrow \Delta} \quad (cut)$$

$\mathcal{P}(C, D)$ :

$$\frac{\Box C \rightarrow \Box C \preceq \Box D, \Box D \prec \Box C \quad \Box D \prec \Box C \rightarrow \Box D \preceq \Box C}{\Box C \rightarrow \Box C \preceq \Box D, \Box D \prec \Box C} \quad (cut) \quad \frac{\Box C \rightarrow \Box C \preceq \Box D, \Box D \prec \Box C}{\Box C \preceq \Box C \rightarrow \Box C \preceq \Box D, \Box D \prec \Box C} \quad (\preceq \rightarrow)$$

$\mathcal{Q}(C, D)$ :

$$\frac{\Box D \rightarrow \Box C \preceq \Box D, \Box D \prec \Box C \quad \Box D \prec \Box C, \Gamma \rightarrow \Delta}{\Box D, \Gamma \rightarrow \Delta, \Box C \preceq \Box D} \quad (cut) \quad \frac{\Box D \preceq \Box C, \Gamma \rightarrow \Delta, \Box C \preceq \Box D}{\Box D \preceq \Box C, \Gamma \rightarrow \Delta, \Box C \preceq \Box D} \quad (\preceq \rightarrow)$$

**Definition 3.3.** A sequent  $\Gamma \rightarrow \Delta$  is said to be saturated if the following conditions hold:

- (1) if  $A \wedge B \in \Gamma$ , then  $A, B \in \Gamma$ , (2) if  $A \wedge B \in \Delta$ , then  $A \in \Delta$  or  $B \in \Delta$ ,
- (3) if  $A \vee B \in \Gamma$ , then  $A \in \Gamma$  or  $B \in \Gamma$ , (4) if  $A \vee B \in \Delta$ , then  $A, B \in \Delta$ ,
- (5) if  $A \supset B \in \Gamma$ , then  $A \in \Delta$  or  $B \in \Gamma$ , (6) if  $A \supset B \in \Delta$ , then  $A \in \Gamma$  and  $B \in \Delta$ ,
- (7) if  $\Box A \preceq \Box B \in \Gamma$ , then  $\Box A \in \Gamma$ , (8) if  $\Box A \preceq \Box A \in \Delta$ , then  $\Box A \in \Delta$ ,
- (9) if  $\Box A \prec \Box B \in \Gamma$ , then  $\Box A \preceq \Box B \in \Gamma$  and  $\Box B \preceq \Box A \in \Delta$ ,
- (10) if  $\Box C, \Box D$  and  $\Box E$  is distinct subformulas in  $\text{Sub}(\Gamma \rightarrow \Delta)$ , then either one of  $\Box C \preceq \Box E \in \Gamma$ ,  $\Box C \preceq \Box D \in \Delta$ , or  $\Box D \preceq \Box E \in \Delta$  holds,
- (11) if  $\Box C$  and  $\Box D$  is distinct subformulas in  $\text{Sub}(\Gamma \rightarrow \Delta)$ , then either one of  $\Box C, \Box D \in \Delta$ ,  $\Box C \prec \Box D \in \Gamma$ ,  $\Box D \prec \Box C \in \Gamma$  or  $\Box C \preceq \Box D, \Box D \preceq \Box C \in \Gamma$  holds.

**Lemma 3.4.** If  $\Gamma \rightarrow \Delta \notin \mathbf{GR}_3^-$ , then there exists a sequent  $\Gamma' \rightarrow \Delta'$  satisfying the following three conditions:

- (1)  $\Gamma' \rightarrow \Delta' \notin \mathbf{GR}_3^-$ ,
- (2)  $\Gamma' \rightarrow \Delta'$  is saturated,
- (3)  $\Gamma \subseteq \Gamma' \subseteq \text{Sub}^+(\Gamma \rightarrow \Delta)$  and  $\Delta \subseteq \Delta' \subseteq \text{Sub}^+(\Gamma \rightarrow \Delta)$ .

Proof. Let it be that  $p \notin \text{Sub}(\Gamma \rightarrow \Delta)$ . Since  $\text{Sub}^+(\Gamma \rightarrow \Delta)$  is finite, there exist formulas

$$A_0, A_1 \cdots, A_{n-1}$$

such that

$$\begin{aligned} \text{Sub}^+(\Gamma \rightarrow \Delta) \cup \{\Box B \wedge \Box C \wedge \Box D \wedge p \mid \Box B, \Box C, \Box D \in \text{Sub}(\Gamma \rightarrow \Delta), B \neq C, C \neq D, D \neq B\} \\ = \{A_0, A_1 \cdots, A_{n-1}\}. \end{aligned}$$

We define a sequence of sequents

$$(\Gamma_0 \rightarrow \Delta_0), (\Gamma_1 \rightarrow \Delta_1), \cdots, (\Gamma_k \rightarrow \Delta_k), \cdots$$

inductively as follows.

**Step 0:**  $(\Gamma_0 \rightarrow \Delta_0) = (\Gamma \rightarrow \Delta)$ .

**Step  $k+1$ :** If  $A_{k \bmod n} = \Box B \preceq \Box C$ , then

$$(\Gamma_{k+1} \rightarrow \Delta_{k+1}) = \begin{cases} (\Box B, \Gamma_k \rightarrow \Delta_k) & \text{if } \Box B \preceq \Box C \in \Gamma \\ (\Gamma_k \rightarrow \Delta_k, \Box B) & \text{if } \Box B \preceq \Box C \in \Delta \text{ and } B = C \\ (\Gamma_k \rightarrow \Delta_k) & \text{otherwise} \end{cases}$$

If  $A_{k \bmod n} = \Box B \prec \Box C$  and  $B \neq C$ , then

$$(\Gamma_{k+1} \rightarrow \Delta_{k+1}) = \begin{cases} S_1 & \text{if } S_1 \notin \mathbf{GR}_3^- \\ (\Box B \preceq \Box C, \Box B \prec \Box C, \Gamma_k \rightarrow \Delta_k, \Box C \preceq \Box B) & \text{if } S_1 \in \mathbf{GR}_3^- \text{ and } S_2 \notin \mathbf{GR}_3^- \\ S_3 & \text{if } S_1, S_2 \in \mathbf{GR}_3^- \text{ and } S_3 \notin \mathbf{GR}_3^- \\ S_4 & \text{if } S_1, S_2, S_3 \in \mathbf{GR}_3^- \text{ and } S_4 \notin \mathbf{GR}_3^- \\ (\Gamma_k \rightarrow \Delta_k) & \text{otherwise} \end{cases}$$

where

$$\begin{aligned} S_1 &= (\Gamma_k \rightarrow \Delta_k, \Box B, \Box C), \\ S_2 &= (\Box B \prec \Box C, \Gamma_k \rightarrow \Delta_k), \\ S_3 &= (\Box C \prec \Box B, \Gamma_k \rightarrow \Delta_k) \text{ and} \\ S_4 &= (\Box B \preceq \Box C, \Box C \preceq \Box B, \Gamma_k \rightarrow \Delta_k). \end{aligned}$$

If  $A_{k \bmod n} = \Box B \prec \Box B$ , then

$$(\Gamma_{k+1} \rightarrow \Delta_{k+1}) = (\Gamma_k \rightarrow \Delta_k)$$

If  $A_{k \bmod n} = \Box B \wedge \Box C \wedge \Box D \wedge p$ , then

$$(\Gamma_{k+1} \rightarrow \Delta_{k+1}) = \begin{cases} S_1 & \text{if } S_1 \notin \mathbf{GR}_1^- \\ S_2 & \text{if } S_1 \in \mathbf{GR}_1^- \text{ and } S_2 \notin \mathbf{GR}_1^- \\ S_3 & \text{if } S_1, S_2 \in \mathbf{GR}_1^- \text{ and } S_3 \notin \mathbf{GR}_1^- \\ (\Gamma_k \rightarrow \Delta_k) & \text{otherwise} \end{cases}$$

where

$$\begin{aligned} S_1 &= (\Gamma_k \rightarrow \Delta_k, \Box B \preceq \Box C), \\ S_2 &= (\Gamma_k \rightarrow \Delta_k, \Box C \preceq \Box D) \text{ and} \\ S_3 &= (\Box B \preceq \Box D, \Gamma_k \rightarrow \Delta_k). \end{aligned}$$

If  $A_{k \bmod n}$  is a  $\Box$ -formula, then  $(\Gamma_{k+1} \rightarrow \Delta_{k+1}) = (\Gamma_k \rightarrow \Delta_k)$ . In the other cases,  $(\Gamma_{k+1} \rightarrow \Delta_{k+1})$  is defined in the usual way.

Also in the usual way, we can prove that  $\bigcup_{i=1}^{\infty} \Gamma_i \rightarrow \bigcup_{i=1}^{\infty} \Delta_i$  is a sequent and satisfies the conditions (1), (2) and (3).  $\dashv$

**Definition 3.5.** For a sequent  $S \notin \mathbf{GR}_3^-$ , we fix a sequent satisfying the three conditions in the above lemma and call it a *saturation* of  $S$ , write  $\text{sat}(S)$ . For  $S \in \mathbf{GR}_3^-$ , we put  $\text{sat}(S) = S$ .

**Remark 3.6.** For a sequent  $S \notin \mathbf{GR}_3^-$ ,

- (1)  $\text{sat}(S) \notin \mathbf{GR}_3^-$ ,
- (2)  $\text{sat}(S)$  is saturated,
- (3)  $\text{Sub}(\text{sat}(S)) = \text{Sub}(S)$ .

A sequence of formulas is defined as follows:

- (1)  $[ ]$  is a sequence of formulas,
- (2) if  $[A_1, \dots, A_n]$  is a sequence of formulas, then so is  $[A_1, \dots, A_n, B]$ .

A binary operator  $\circ$  is defined by

$$[A_1, \dots, A_m] \circ [B_1, \dots, B_n] = [A_1, \dots, A_m, B_1, \dots, B_n].$$

We use  $\tau$  and  $\sigma$ , possibly with suffixes, for sequences of formulas.

**Definition 3.7.** Let  $S_0$  be a sequent, which is not provable in  $\mathbf{GR}_3^-$ . We define the set  $\mathbf{W}(S_0)$  of pairs of a sequent and a sequence of formulas as follows:

- (1)  $(\text{sat}(S_0); [ ]) \in \mathbf{W}(S_0)$ ,
- (2) if a pair  $(\Gamma \rightarrow \Delta, \Box A; \tau)$  belongs to  $\mathbf{W}(S_0)$ , then so does the pair

$$(\text{sat}(\Box A, \{D \mid \Box D \in \Gamma\}, \Gamma \cap \Sigma \rightarrow A); \tau \circ [\Box A]).$$

**Lemma 3.8.** Let  $S_0$  be a sequent, which is not provable in  $\mathbf{GR}_3^-$  and let  $(S; \tau)$  be a pair in  $\mathbf{W}(S_0)$ . Then

- (1)  $S$  is saturated,
- (2)  $S \notin \mathbf{GR}_3^-$ ,
- (3)  $S$  consists of only formulas in  $\text{Sub}^+(S_0)$ ,
- (4)  $\tau$  consists of only  $\Box$ -formulas in  $\text{Sub}(S_0)$ .

*Proof.* We use an induction on  $(S; \tau)$  as an element in  $\mathbf{W}(S_0)$ . If  $(S; \tau) = (\text{sat}(S_0); [ ])$ , then the lemma is clear from Definition 3.5. Suppose that  $(S; \tau) \neq (\text{sat}(S_0); [ ])$ . Then by Definition 3.7, there exists a pair  $(\Gamma \rightarrow \Delta, \Box A; \sigma) \in \mathbf{W}(S_0)$  such that

$$(S; \tau) = (\text{sat}(\Box A, \{D \mid \Box D \in \Gamma\}, \Gamma \cap \Sigma \rightarrow A); \sigma \circ [\Box A]).$$

So, we obtain (1). By the induction hypothesis, we have the following three:

- (5)  $\Gamma \rightarrow \Delta, \Box A \notin \mathbf{GR}_3^-$ ,
- (6)  $\Gamma \rightarrow \Delta, \Box A$  consists only formulas in  $\text{Sub}^+(S_0)$



(7)  $\sigma$  consists of only  $\Box$ -formulas in  $\text{Sub}(S_0)$ .

From (6) and (7), we obtain (4). By Remark 3.6(3) and (6), we have (3). Also by (5), we have  $\Gamma \cap \Sigma \rightarrow \Box A \notin \mathbf{GR}_3^-$ . Using  $(\rightarrow \Box)$ , we have  $\Box A, \{D \mid \Box D \in \Gamma\}, \Gamma \cap \Sigma \rightarrow A \notin \mathbf{GR}_3^-$ , and by Remark 3.6(1), neither is its saturation  $S$ . We have (2).  $\dashv$

**Lemma 3.9.** *Let  $S_0$  be a sequent, which is not provable in  $\mathbf{GR}_3^-$ . Then*

- (1)  $S_1 = S_2$  for any  $(S_1; \tau), (S_2; \tau) \in \mathbf{W}(S_0)$ ,
- (2)  $(\Gamma_1 \rightarrow \Delta_1; \tau), (\Gamma_2 \rightarrow \Delta_2; \tau \circ \sigma) \in \mathbf{W}(S_0)$  implies  $\Gamma_1 \cap \Sigma \subseteq \Gamma_2$ ,
- (3) if there exists a  $\Box$ -formula in the antecedent of  $\text{sat}(S_0)$ , then  $\text{Sub}^+(S_0) \cap \Sigma = \text{Sub}^+(S) \cap \Sigma$ , for any  $(S; \tau) \in \mathbf{W}(S_0)$ ,
- (4)  $(\Gamma \rightarrow \Delta; \tau \circ \sigma) \in \mathbf{W}(S_0)$  implies  $(\Gamma_1 \rightarrow \Delta_1; \tau) \in \mathbf{W}(S_0)$  for some  $\Gamma_1 \rightarrow \Delta_1$ ,
- (5)  $(\Gamma \rightarrow \Delta; \tau \circ [\Box A] \circ \sigma) \in \mathbf{W}(S_0)$  implies  $\Box A \in \Gamma$ ,
- (6)  $(S; \tau_1 \circ [\Box A] \circ \tau_2 \circ [\Box A] \circ \tau_3) \notin \mathbf{W}(S_0)$ , for any  $A$  and  $S$ ,
- (7)  $\mathbf{W}(S_0)$  is finite.

Proof. For (1). We use an induction on  $\tau$ . If  $\tau = [ ]$ , then we have  $S_1 = S_2 = \text{sat}(S_0)$ . Suppose that  $\tau = \sigma \circ [\Box A]$ . Then by Definition 3.7, there exist

$$(\Gamma_1 \rightarrow \Delta_1, \Box A; \sigma), (\Gamma_2 \rightarrow \Delta_2, \Box A; \sigma) \in \mathbf{W}(S_0)$$

such that

$$S_1 = \text{sat}(\Box A, \{D \mid \Box D \in \Gamma_1\}, \Gamma_1 \cap \Sigma \rightarrow A)$$

and

$$S_2 = \text{sat}(\Box A, \{D \mid \Box D \in \Gamma_2\}, \Gamma_2 \cap \Sigma \rightarrow A)$$

By the induction hypothesis, we have  $(\Gamma_1 \rightarrow \Delta_1) = (\Gamma_2 \rightarrow \Delta_2)$ , and so  $\Gamma_1 = \Gamma_2$ . Hence we obtain  $S_1 = S_2$ .

For (2). We use an induction on  $\sigma$ . If  $\sigma = [ ]$ , then by (1), we have  $\Gamma_1 = \Gamma_2$ . Suppose that  $\sigma = \sigma' \circ [\Box A]$ . Then by Definition 3.7, there exists  $(\Gamma_3 \rightarrow \Delta_3, \Box A; \tau \circ \sigma') \in \mathbf{W}(S_0)$  such that  $(\Gamma_2 \rightarrow \Delta_2) = \text{sat}(\Box A, \{D \mid \Box D \in \Gamma_3\}, \Gamma_3 \cap \Sigma \rightarrow A)$ . By the induction hypothesis,  $\Gamma_1 \cap \Sigma \subseteq \Gamma_3$ . Using Remark 3.6(3),  $\Gamma_1 \cap \Sigma \subseteq \Gamma_3 \cap \Sigma \subseteq \Gamma_2$ .

For (3). By Lemma 3.8(2), we have  $\text{Sub}^+(S_0) \cap \Sigma \supseteq \text{Sub}^+(S) \cap \Sigma$ . We show  $\text{Sub}^+(S_0) \cap \Sigma \subseteq \text{Sub}^+(S) \cap \Sigma$ . We put  $S_0 = \Gamma_0 \rightarrow \Delta_0$ . Suppose that  $\Box C \in \Gamma_0$  and  $E \in \text{Sub}^+(S_0) \cap \Sigma$ .

If  $E = \Box A$ , then either one of the following four holds since  $\text{sat}(S_0)$  is saturated:

- (3a)  $\Box C, \Box A \in \Delta_0$ ,
- (3b)  $\Box C \prec \Box A, \Box C \preceq \Box A \in \Gamma_0$ ,
- (3c)  $\Box A \prec \Box C, \Box A \preceq \Box C \in \Gamma_0$ ,
- (3d)  $\Box C \preceq \Box A, \Box A \preceq \Box C \in \Gamma_0$ .

By Lemma 3.8(1) and  $\Box C \in \Gamma_0$ , we note that (3a) does not hold. So, one of the formulas  $\Box C \prec \Box A$  and  $\Box A \preceq \Box C$  belongs to  $\Gamma_0$ , and by (2), it also belongs to the antecedent of  $S$ . Hence  $\Box A \in \text{Sub}^+(S) \cap \Sigma$ .

If  $E$  is either a  $\preceq$ -formula  $\Box A \preceq \Box B$  or a  $\prec$ -formula  $\Box A \prec \Box B$ , then  $\Box A, \Box B \in \text{Sub}^+(S_0) \cap \Sigma$ . So similarly to the proof for the case that  $E = \Box A$ , we have  $\Box A, \Box B \in \text{Sub}^+(S) \cap \Sigma$ , and so,  $E \in \text{Sub}^+(S) \cap \Sigma$ .

For (4). We use an induction on  $\sigma$ . If  $\sigma = [ ]$ , then the lemma is clear. Suppose that  $\sigma = \sigma' \circ [A]$ . Then by Definition 3.7, there exists  $(\Gamma_1 \rightarrow \Delta_1, \Box A; \tau \circ \sigma') \in \mathbf{W}(S_0)$  such that

$$(\Gamma \rightarrow \Delta) = \text{sat}(\Box A, \{D \mid \Box D \in \Gamma_1\}, \Gamma_1 \cap \Sigma \rightarrow A).$$

By the induction hypothesis,  $(\Gamma_2 \rightarrow \Delta_2; \tau) \in \mathbf{W}(S_0)$  for some  $\Gamma_2 \rightarrow \Delta_2$ .

For (5). By (4),  $(\Gamma_1 \rightarrow \Delta_1; \tau \circ [\Box A]) \in \mathbf{W}(S_0)$  for some  $\Gamma_1 \rightarrow \Delta_1$ . Using Definition 3.7, we have  $\Box A \in \Gamma_1$ . Using (2), we have  $\Box A \in \Gamma$ .

For (6). Suppose that  $(S; \tau_1 \circ [\Box A] \circ \tau_2 \circ [\Box A] \circ \tau_3) \in \mathbf{W}(S_0)$ . Then by (4),  $(\Gamma \rightarrow \Delta; \tau_1 \circ [\Box A] \circ \tau_2 \circ [\Box A]) \in \mathbf{W}(S_0)$  for some  $\Gamma \rightarrow \Delta$ . By Definition 3.7, there exists  $(\Gamma_1 \rightarrow \Delta_1, \Box A; \tau_1 \circ [\Box A] \circ \tau_2) \in \mathbf{W}(S_0)$ . Using (5),  $\Box A \in \Gamma_1$ . So,  $\Gamma_1 \rightarrow \Delta_1, \Box A \in \mathbf{GR}_3^-$ . This is contradictory to Lemma 3.8.

For (7). By (1), (6) and Lemma 3.8(4).  $\dashv$

**Definition 3.10.** Let  $S_0$  be a sequent, which is not provable in  $\mathbf{GR}_3^-$ . We define a structure  $\mathcal{K}(S_0) = \langle \mathbf{W}(S_0), <, \models \rangle$  as follows:

- (1)  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) < (\Gamma_2 \rightarrow \Delta_2; \tau_2)$  if and only if  $\tau_2 = \tau_1 \circ \sigma$  for some  $\sigma \neq [\ ]$ ,
- (2)  $\models$  is a valuation satisfying, in addition to the conditions in Definition 1.1(3),
  - (2.1)  $p \in \Gamma$  if and only if  $(\Gamma \rightarrow \Delta; \tau) \models p$ , for any propositional variable  $p$ ,
  - (2.2)  $A \in \Gamma$  if and only if  $(\Gamma \rightarrow \Delta; \tau) \models A$ , for any  $\prec$ -formula  $A \in \text{Sub}^+(S_0)$ ,
  - (2.3)  $(\Gamma \rightarrow \Delta; \tau) \models \Box A$  if and only if  $(\Gamma \rightarrow \Delta; \tau) \models \Box A \preceq \Box A$ , for any  $\Box A \in \text{Sub}(S_0)$ ,
  - (2.4)  $\Box A \preceq \Box B \in \Gamma$  if and only if  $(\Gamma \rightarrow \Delta; \tau) \models \Box A \preceq \Box B$ , for any  $\Box A, \Box B \in \text{Sub}(S_0)$  such that  $A \neq B$ .

**Lemma 3.11.** Let  $S_0$  be a sequent, which is not provable in  $\mathbf{GR}_3^-$ . Then for any  $A$  and for any  $(\Gamma \rightarrow \Delta; \tau) \in \mathbf{W}(S_0)$ ,

- (1)  $A \in \Gamma$  implies  $(\Gamma \rightarrow \Delta; \tau) \models A$ ,
- (2)  $A \in \Delta$  implies  $(\Gamma \rightarrow \Delta; \tau) \not\models A$ .

*Proof.* We use an induction on  $A$ .

If  $A = \perp$ , then by Lemma 3.8(1),  $A \notin \Gamma$ . So we have (1). On the other hand, from  $(\Gamma \rightarrow \Delta; \tau) \not\models A$ , we have (2).

If  $A$  is a propositional variable, then (1) is clear. Suppose that  $p \in \Delta$ . By Lemma 3.8(1),  $p \notin \Gamma$ , and so, we have (2).

Suppose that  $A$  is not a propositional variable. If  $A$  is a  $\preceq$ -formula  $\Box B \preceq \Box C$  with  $B \neq C$  or a  $\prec$ -formula, then the lemma can be shown similarly to the case that  $A$  is a propositional variable. Other cases can be shown in the usual way (cf. [Avr84]). Here we show the case that  $A = \Box B \preceq \Box B$  and the case that  $A = \Box B$ .

For the case that  $A = \Box B \preceq \Box B$ . Suppose that  $\Box B \preceq \Box B \in \Gamma$ . Since  $\Gamma \rightarrow \Delta$  is saturated, we have  $\Box B \in \Gamma$ . By the induction hypothesis,  $(\Gamma \rightarrow \Delta; \tau) \models \Box B$ . From Definition 2.8(2.4), we obtain  $(\Gamma \rightarrow \Delta; \tau) \models \Box B \preceq \Box B$ .

Suppose that  $\Box B \preceq \Box B \in \Delta$ . Since  $\Gamma \rightarrow \Delta$  is saturated, we have  $\Box B \in \Delta$ . By the induction hypothesis,  $(\Gamma \rightarrow \Delta; \tau) \not\models \Box B$ . From Definition 2.8(2.4), we obtain  $(\Gamma \rightarrow \Delta; \tau) \not\models \Box B \preceq \Box B$ .

For the case that  $A = \Box B$ . Suppose that  $\Box B \in \Gamma$  and  $(\Gamma \rightarrow \Delta; \tau) < (\Gamma_1 \rightarrow \Delta_1; \tau_1)$ . Then  $\tau_1 = \tau \circ \sigma \circ [\Box C]$  for some  $\sigma$  and  $C$ . Hence there exists  $(\Gamma_2 \rightarrow \Delta_2, \Box C; \tau \circ \sigma) \in \mathbf{W}(S_0)$  such that

$$(\Gamma_1 \rightarrow \Delta_1) = \text{sat}(\Box C, \{D \mid \Box D \in \Gamma_2\}, \Gamma_2 \cap \Sigma \rightarrow C).$$

By Lemma 3.9(2), we have  $\Box B \in \Gamma_2$ . Using Definition 3.5,  $B \in \Gamma_1$ . By the induction hypothesis, we have  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \models B$ . Hence  $(\Gamma \rightarrow \Delta; \tau) \models \Box B$ .

Suppose that  $\Box B \in \Delta$ . Then  $(\Gamma \rightarrow \Delta; \tau) < (\text{sat}(\Box B, \{D \mid \Box D \in \Gamma\}, \Gamma \cap \Sigma \rightarrow B); \tau \circ [\Box B]) \in \mathbf{W}(S_0)$ . By Definition 3.5,  $B$  belongs to the succedent of the above saturation. By the induction hypothesis,  $B$  is false at the new pair above. Hence  $(\Gamma \rightarrow \Delta; \tau) \not\models \Box B$ .  $\dashv$

**Corollary 3.12.** Let  $A_1, \dots, A_m \rightarrow B_1, \dots, B_n$  be a sequent, which is not provable in  $\mathbf{GR}_3^-$ . Then in  $\mathcal{K}(A_1, \dots, A_m \rightarrow B_1, \dots, B_n)$ ,

$$(\text{sat}(A_1, \dots, A_m \rightarrow B_1, \dots, B_n); [\ ]) \not\models A_1 \wedge \dots \wedge A_m \supset B_1 \vee \dots \vee B_n.$$

**Lemma 3.13.** Let  $S_0$  be a sequent, which is not provable in  $\mathbf{GR}_3^-$ . If there exists a  $\Box$ -formula in the antecedent of  $\text{sat}(S_0)$ , then  $\mathcal{K}(S_0)$  is a Kripke pseudo-model for  $\mathbf{R}^-$  satisfying the two conditions in Definition 1.2 for any  $D$  such that  $\text{Sub}(D) \cap \Sigma \subseteq \text{Sub}^+(S_0)$ .

*Proof.* By Lemma 3.9(7),  $\mathbf{W}(S_0)$  is finite. The irreflexivity and the transitivity of  $<$  can be shown easily. We show

$$\alpha < \gamma \text{ and } \beta < \gamma \text{ imply either one of } \alpha = \beta, \alpha < \beta \text{ or } \beta < \alpha.$$

Suppose that  $(S_1; \tau_1) < (S_3; \tau_3)$  and  $(S_2; \tau_2) < (S_3; \tau_3)$ . Then  $\tau_3 = \tau_1 \circ \sigma_1 = \tau_2 \circ \sigma_2$  for some non-empty sequences  $\sigma_1$  and  $\sigma_2$ . Hence either  $\tau_1 = \tau_2 \circ \sigma'_2$  or  $\tau_1 \circ \sigma'_1 = \tau_2$  holds. Using Lemma 3.9(1), we have either one of  $(S_1; \tau_1) = (S_2; \tau_2)$ ,  $(S_1; \tau_1) < (S_2; \tau_2)$  or  $(S_2; \tau_2) < (S_1; \tau_1)$ . Hence  $\mathcal{K}(S_0)$  is a Kripke pseudo-model for  $\mathbf{R}^-$ .

We show the two conditions in Definition 1.2 for any  $D$  such that  $\text{Sub}(D) \cap \Sigma \subseteq \text{Sub}^+(S_0)$ .

For (1). If  $D$  is  $\Box$ -formula, then (1) is clear by the definition of  $\models$ . Suppose that  $D$  is either a  $\preceq$ -formula or a  $\prec$ -formula,  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \models D$  and  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) < (\Gamma_2 \rightarrow \Delta_2; \tau_2)$ . Then  $D \in \Gamma_1$ . By Lemma 3.9(2),  $D \in \Gamma_2$ . Hence  $(\Gamma_2 \rightarrow \Delta_2; \tau_2) \models D$ .

For (2). We divide the cases.

For the case that  $D$  is A5 (i.e.,  $D = (\Box A \preceq \Box B) \supset \Box A$ ). Suppose that  $(\Gamma \rightarrow \Delta; \tau) \models \Box A \preceq \Box B$ . If  $A = B$ , then immediately  $(\Gamma \rightarrow \Delta; \tau) \models \Box A$ . So, we assume that  $A \neq B$ . Since  $\text{Sub}(D) \cap \Sigma \subseteq \text{Sub}^+(S_0)$ , we have  $\Box A \preceq \Box B \in \text{Sub}^+(S_0)$ , and so  $\Box A \preceq \Box B \in \Gamma$ . Since  $\Gamma \rightarrow \Delta$  is saturated, we have  $\Box A \in \Gamma$ . Using Lemma 3.11,  $(\Gamma \rightarrow \Delta; \tau) \models \Box A$ .

For the case that  $D$  is A6 (i.e.,  $D = ((\Box A \preceq \Box B) \wedge (\Box B \preceq \Box C)) \supset (\Box A \preceq \Box C)$ ). Suppose that  $(\Gamma \rightarrow \Delta; \tau) \models \Box A \preceq \Box B$  and  $(\Gamma \rightarrow \Delta; \tau) \models \Box B \preceq \Box C$ . If either  $A = B$  or  $B = C$ , then  $(\Gamma \rightarrow \Delta; \tau) \models \Box A \preceq \Box C$  is clear. So, we assume that  $A \neq B$  and  $B \neq C$ . Since  $\text{Sub}(D) \cap \Sigma \subseteq \text{Sub}^+(S_0)$ , we have  $\Box A \preceq \Box B, \Box B \preceq \Box C \in \text{Sub}^+(S_0)$ , and so,  $\Box A \preceq \Box B, \Box B \preceq \Box C \in \Gamma$ .

If  $A \neq C$ , then  $\Box A \preceq \Box C \in \Gamma$  since  $\Gamma \rightarrow \Delta$  is saturated. Using Lemma 3.11,  $(\Gamma \rightarrow \Delta; \tau) \models \Box A \preceq \Box C$ .

If  $A = C$ , then  $\Box A \in \Gamma$  since  $\Gamma \rightarrow \Delta$  is saturated. Using Lemma 3.11,  $(\Gamma \rightarrow \Delta; \tau) \models \Box A$ , and so  $(\Gamma \rightarrow \Delta; \tau) \models \Box A \preceq \Box C$ .

For the case that  $D$  is A7 (i.e.,  $D = (\Box A \vee \Box B) \supset ((\Box A \preceq \Box B) \vee (\Box B \prec \Box A))$ ). Suppose that  $(\Gamma \rightarrow \Delta; \tau) \models \Box A \vee \Box B$ , i. e., either  $\Box A$  or  $\Box B$  is true at the pair.

If  $A = B$ , then  $\Box A$  is true at  $(\Gamma \rightarrow \Delta; \tau)$ , and so are  $\Box A \preceq \Box B$  and  $(\Box A \preceq \Box B) \vee (\Box B \prec \Box A)$ .

So, we assume that  $A \neq B$ . Since  $\text{Sub}(D) \cap \Sigma \subseteq \text{Sub}^+(S_0)$ , we have  $\Box A, \Box B \in \text{Sub}^+(S_0)$ , and using Lemma 3.9(3), we also have  $\Box A, \Box B \in \text{Sub}^+(\Gamma \rightarrow \Delta)$ . Since  $\Gamma \rightarrow \Delta$  is saturated, either one of the following four conditions holds:

- $\Box A, \Box B \in \Delta$ ,
- $\Box A \prec \Box B, \Box A \preceq \Box B \in \Gamma$ ,
- $\Box B \prec \Box A, \Box B \preceq \Box A \in \Gamma$ ,
- $\Box A \preceq \Box B, \Box B \preceq \Box A \in \Gamma$ .

If  $\Box A, \Box B \in \Delta$ , then by Lemma 3.11,  $\Box A$  and  $\Box B$  are false at  $(\Gamma \rightarrow \Delta; \tau)$ , which is in contradiction with  $(\Gamma \rightarrow \Delta; \tau) \models \Box A \vee \Box B$ . So, either one of  $\Box A \preceq \Box B$  or  $\Box B \prec \Box A$  is true at  $(\Gamma \rightarrow \Delta; \tau)$ , and so is  $(\Box A \preceq \Box B) \vee (\Box B \prec \Box A)$ .

For the case that  $D$  is A8 (i.e.,  $D = (\Box A \prec \Box B) \supset (\Box A \preceq \Box B)$ ). Suppose that  $(\Gamma \rightarrow \Delta; \tau) \models \Box A \prec \Box B$ . Since  $\text{Sub}(D) \cap \Sigma \subseteq \text{Sub}^+(S_0)$ , we have  $\Box A \prec \Box B \in \text{Sub}^+(S_0)$ , and so,  $\Box A \prec \Box B \in \Gamma$ . Since  $\Gamma \rightarrow \Delta$  is saturated,  $\Box A \preceq \Box B \in \Gamma$ , and by Lemma 3.11,  $(\Gamma \rightarrow \Delta; \tau) \models \Box A \preceq \Box B$ .

For the case that  $D$  is A9 (i.e.,  $D = ((\Box A \preceq \Box B) \wedge (\Box B \prec \Box A)) \supset \perp$ ). Suppose that  $(\Gamma \rightarrow \Delta; \tau) \models \Box A \preceq \Box B$  and  $(\Gamma \rightarrow \Delta; \tau) \models \Box B \prec \Box A$ . Since  $\text{Sub}(D) \cap \Sigma \subseteq \text{Sub}^+(S_0)$ , we have  $\Box B \prec \Box A \in \Gamma$ . Since  $\Gamma \rightarrow \Delta$  is saturated,  $\Box A \preceq \Box B \in \Delta$ . and by Lemma 3.11,  $(\Gamma \rightarrow \Delta; \tau) \models \Box A \preceq \Box B$ , which is contradictory to  $(\Gamma \rightarrow \Delta; \tau) \models \Box B \prec \Box A$ .  $\dashv$

**Lemma 3.14**[(GS79)]. *Let  $\mathbf{S}$  be a set of formulas satisfying*

$$A \in \mathbf{S} \text{ implies } \text{Sub}^+(A) \subseteq \mathbf{S}$$

*and Let  $\mathcal{K}^*$  be a Kripke pseudo-model for  $\mathbf{R}^-$  satisfying the two conditions in Definition 1.2 for any  $D$  such that  $\text{Sub}(D) \cap \Sigma \subseteq \mathbf{S}$ . Then there exists a Kripke model  $\mathcal{K}$  for  $\mathbf{R}^-$  such that for any  $A \in \mathbf{S}$ ,*

$$A \text{ is valid in } \mathcal{K}^* \text{ if and only if } A \text{ is valid in } \mathcal{K}$$

The lemma above sometimes called “extension lemma”.

**Theorem 3.15.** *Let  $A_1, \dots, A_m \rightarrow B_1, \dots, B_n$  be a sequent, which is not provable in  $\mathbf{GR}_3^-$ . Then there exists a Kripke model  $\mathbf{K}$  for  $\mathbf{R}^-$ , in which the formula  $A_1 \wedge \dots \wedge A_m \supset B_1 \vee \dots \vee B_n$  is not valid.*

Proof. Let  $\Gamma$  and  $\Delta$  be the antecedent and succedent of  $\text{sat}(A_1, \dots, A_m \rightarrow B_1, \dots, B_n)$ , respectively. For any  $\Box A \in \Delta$ ,  $S(A)$  denotes the sequent  $\Box A, \{D \mid \Box D \in \Gamma\}, \Gamma \cap \Sigma \rightarrow A$ . By  $(\rightarrow \Box)$ , we note  $S(A) \notin \mathbf{GR}_3^-$ . Also we note that there exists a  $\Box$ -formula in the antecedent of the saturation of  $S(A)$ . Using Corollary 3.12, Lemma 3.13 and Lemma 3.14, there exists a Kripke model  $\mathcal{K}_A = \langle \mathbf{W}_A, <_A, \models_A \rangle$  for  $\mathbf{R}^-$  such that for any  $B \in \text{Sub}^+(S(A))$ ,

$B$  is valid in  $\mathcal{K}(S(A))$  if and only if  $B$  is valid in  $\mathcal{K}_A$ .

We construct a structure  $\mathcal{K} = \langle \mathbf{W}, <, \models \rangle$  as follows:

- (1)  $\mathbf{W} = \{r\} \cup \bigcup_{\Box A \in \Delta} \{(w, A) \mid w \in \mathbf{W}_A\}$ ,
- (2)  $<$  is a binary relation on  $\mathbf{W}$  satisfying
  - (2.1)  $r < \alpha$  if and only if  $\alpha \neq r$ ,
  - (2.2)  $(w_1, A_1) < (w_2, A_2)$  if and only if  $A_1 = A_2$  and  $w_1 <_{A_1} w_2$ ,
- (3)  $\models$  is a valuation satisfying, in addition to the conditions in Definition 1.1(3),
  - (3.1)  $r \models \Box B \preceq \Box B$  if and only if  $r \models \Box B$ ,
  - (3.2)  $(w, A) \models \Box B \preceq \Box B$  if and only if  $w \models_A \Box B$ ,
  - (3.3)  $r \models C$  if and only if  $C \in \Gamma$ ,
  - (3.4)  $(w, A) \models C$  if and only if  $w \models_A C$ ,

where  $C$  is either one of a propositional variable, a  $\prec$ -formula  $\Box B \preceq \Box C$  with  $B \neq C$  or a  $\preceq$ -formula. We note that for any formula  $D$  and for any  $w \in \mathbf{W}_A$ ,

$$w \models_A D \text{ if and only if } (w, A) \models D. \dots \dots (*)$$

We show

- (4)  $B \in \Gamma$  implies  $r \models B$

and

- (5)  $B \in \Delta$  implies  $r \not\models B$

by an induction on  $B$ .

If  $B$  is either one of a propositional variable, a  $\prec$ -formula  $\Box C \preceq \Box D$  with  $C \neq D$  or a  $\preceq$ -formula, then (4) is clear. Suppose that  $B \in \Delta$ . Then we have  $B \notin \Gamma$  since  $\Gamma \rightarrow \Delta \notin \mathbf{GR}_3^-$ . By (3.3), we have  $r \not\models B$ .

Among the other cases, we only show the case that  $B = \Box C \preceq \Box C$  and the case that  $B = \Box C$ .

For the case that  $B = \Box C \preceq \Box C$ .

$$\begin{aligned} & \Box C \preceq \Box C \in \Gamma \\ \Rightarrow & \Box C \in \Gamma && \text{since } \Gamma \rightarrow \Delta \text{ is saturated} \\ \Rightarrow & r \models \Box C && \text{by the induction hypothesis} \\ \Rightarrow & r \models \Box C \preceq \Box C && \text{by (3.3)} \end{aligned}$$

and

$$\begin{aligned} & \Box C \preceq \Box C \in \Delta \\ \Rightarrow & \Box C \in \Delta && \text{since } \Gamma \rightarrow \Delta \text{ is saturated} \\ \Rightarrow & r \not\models \Box C && \text{by the induction hypothesis} \\ \Rightarrow & r \not\models \Box C \preceq \Box C && \text{by (3.3).} \end{aligned}$$

For the case that  $B = \Box C$ .

Suppose that  $\Box C \in \Gamma$  and let  $\Box A$  be an formula in  $\Delta$ . Then  $C, \Box C$  belong to the antecedent of  $S(A)$ . By Lemma 3.11, in  $\mathcal{K}(S(A))$ ,

$$(\text{sat}(S(A)); [\ ] \models C \text{ and } (\text{sat}(S(A)); [\ ] \models \Box C.$$

Since  $(\text{sat}(S(A)); [\ ])$  is the root of  $\mathcal{K}(S(A))$ ,  $C$  is valid in  $\mathcal{K}(S(A))$ . Hence  $C$  is valid in  $\mathcal{K}_A$ . Using (\*), in  $\mathcal{K}$ , we have  $(w, A) \models B$  for any  $w \in \mathbf{W}_A$ . Hence  $r \models \Box B$ .

Suppose that  $\Box C \in \Delta$ . Then  $C$  belongs the succedent of  $S(C)$ . By Lemma 3.11, in  $\mathcal{K}(S(C))$ ,

$$(\text{sat}(S(A)); [\ ]) \not\models C.$$

So,  $C$  is not valid in  $\mathcal{K}_B$ , and there exists  $w \in \mathbf{W}_B$  such that  $w \not\models C$ . Using (\*), in  $\mathcal{K}$ ,  $(w, C) \not\models C$ . Since  $r < (w, B)$ , we have  $r \not\models \Box C$ .

Hence we obtain (4) and (5), and so,

$$r \not\models (A_1 \wedge \cdots \wedge A_m) \supset (B_1 \vee \cdots \vee B_n).$$

We show  $\mathcal{K}$  is a pseudo-Kripke model for  $\mathbf{R}^-$  satisfying two conditions in Definition 1.2 for any  $D$  such that  $\mathbf{Sub}(D) \cap \Sigma \subseteq \mathbf{Sub}^+(A_1, \dots, A_m \rightarrow B_1, \dots, B_n)$ . Since  $\mathcal{K}_A$  is a Kripke model, we can see that  $\mathcal{K}$  is a Kripke pseudo-model for  $\mathbf{R}^-$  from the definition of  $\mathcal{K}$ . Also since  $\mathcal{K}_A$  is a Kripke model, two conditions in Definition 1.2 are clear if  $\alpha \neq r$ . So, it is sufficient to show the following two:

- (6) if  $D \in \Sigma$  and  $r \models D$ , then for any  $\Box A \in \Delta$  and for any  $w \in \mathbf{W}_A$ ,  $(w, A) \models D$ ,
- (7) if  $D$  is either one of the axioms  $A5, A6, A7, A8$  and  $A9$ , then  $r \models D$ .

For (6): If  $D = \Box E$ , then (6) is clear by the definition of  $\models$ . If  $D = \Box E \prec \Box F$ , then by (3.1), we obtain (6). So, we assume that  $D = \Box E \preceq \Box F$  with  $E \neq F$  or  $D$  is a  $\prec$ -formula. Suppose that  $r \models D$  and let it be that  $\Box A \in \Delta$ . Then by (3.3),  $D \in \Gamma \cap \Sigma$ . Hence  $D$  belongs the antecedent of  $\text{sat}(S(A))$ . Using Lemma 3.9(2) and  $(\text{sat}(S(A)); [\ ] \in \mathbf{W}(S(A)))$ ,  $D \in \Phi$  for any  $(\Phi \rightarrow \Psi; \sigma) \in \mathbf{W}(S(A))$ . Using Lemma 3.11,  $D$  is valid in  $\mathcal{K}(S(A))$ , and hence,  $D$  is valid in  $\mathcal{K}_A$ . Using  $(*)$ , we obtain for any  $w \in \mathbf{W}_A$ ,  $(w, A) \models D$ .

For (7): We divide the cases.

For the case that  $D$  is  $A5$  (i.e.,  $D = (\Box E \preceq \Box F) \supset \Box E$ ). If  $E = F$ , then by (3.1),

$$r \models \Box E \preceq \Box F \Rightarrow r \models \Box E.$$

If  $E \neq F$ , then

$$\begin{aligned} & r \models \Box E \preceq \Box F \\ \Rightarrow & \Box E \preceq \Box F \in \Gamma && \text{by (3.3)} \\ \Rightarrow & \Box E \in \Gamma && \text{since } \Gamma \rightarrow \Delta \text{ is saturated} \\ \Rightarrow & r \models \Box E && \text{by (4).} \end{aligned}$$

For the case that  $D$  is  $A6$  (i.e.,  $D = ((\Box E \preceq \Box F) \wedge (\Box F \preceq \Box G)) \supset (\Box E \preceq \Box G)$ ). If  $E = F$  or  $F = G$ , then

$$r \models \Box E \preceq \Box F, r \models \Box F \preceq \Box G \Rightarrow r \models \Box E \preceq \Box G.$$

If  $E \neq F$  and  $E = G$ , then

$$\begin{aligned} & r \models \Box E \preceq \Box F, r \models \Box F \preceq \Box E \\ \Rightarrow & \Box E \preceq \Box F \in \Gamma && \text{by (3.3)} \\ \Rightarrow & \Box E \in \Gamma && \text{since } \Gamma \rightarrow \Delta \text{ is saturated} \\ \Rightarrow & r \models \Box E && \text{by (4)} \\ \Rightarrow & r \models \Box E \preceq \Box E && \text{by (3.1).} \end{aligned}$$

If  $E \neq F, F \neq G$  and  $E \neq G$ , then

$$\begin{aligned} & r \models \Box E \preceq \Box F, r \models \Box F \preceq \Box G \\ \Rightarrow & \Box E \preceq \Box F, \Box F \preceq \Box G \in \Gamma && \text{by (3.3)} \\ \Rightarrow & \Box E \prec \Box G \in \Gamma && \text{since } \Gamma \rightarrow \Delta \text{ is saturated} \\ \Rightarrow & r \models \Box E \prec \Box G && \text{by (4).} \end{aligned}$$

For the case that  $D$  is  $A7$  (i.e.,  $D = (\Box E \vee \Box F) \supset ((\Box E \preceq \Box F) \vee (\Box F \prec \Box E))$ ). If  $E = F$ , then

$$\begin{aligned} & r \models \Box E \text{ or } r \models \Box F \\ \Rightarrow & r \models \Box E \preceq \Box E && \text{by (3.1)} \\ \Rightarrow & r \models (\Box E \prec \Box F) \vee (\Box F \prec \Box E). \end{aligned}$$

If  $E \neq F$ , then

$$\begin{aligned} & r \models \Box E \text{ or } r \models \Box F \\ \Rightarrow & \Box E \notin \Delta \text{ or } \Box F \notin \Delta && \text{by (5)} \\ \Rightarrow & \Box E \preceq \Box F \in \Gamma \text{ or } \Box F \preceq \Box E \in \Gamma && \text{since } \Gamma \rightarrow \Delta \text{ is saturated} \\ \Rightarrow & r \models \Box E \preceq \Box F \in \Gamma \text{ or } r \models \Box F \preceq \Box E \in \Gamma && \text{by (4).} \end{aligned}$$

For the case that  $D$  is A8 (i.e.,  $D = (\Box E \prec \Box F) \supset (\Box F \preceq \Box E)$ ).

$$\begin{aligned}
& r \models \Box E \prec \Box F \\
\Rightarrow & \Box E \prec \Box F \in \Gamma && \text{by (3.3)} \\
\Rightarrow & \Box E \preceq \Box F \in \Gamma && \text{since } \Gamma \rightarrow \Delta \text{ is saturated} \\
\Rightarrow & r \models \Box E \preceq \Box F && \text{by (4).}
\end{aligned}$$

For the case that  $D$  is A9 (i.e.,  $D = ((\Box E \preceq \Box F) \wedge (\Box F \prec \Box E)) \supset \perp$ ).

$$\begin{aligned}
& r \models \Box F \prec \Box E \\
\Rightarrow & \Box F \prec \Box E \in \Gamma && \text{by (3.3)} \\
\Rightarrow & \Box E \preceq \Box F \in \Delta && \text{since } \Gamma \rightarrow \Delta \text{ is saturated} \\
\Rightarrow & r \not\models \Box E \preceq \Box F && \text{by (5).}
\end{aligned}$$

Hence using Lemma 3.14, we obtain the theorem.  $\dashv$

**Corollary 3.16.** *The following conditions are equivalent:*

- (1)  $A_1, \dots, A_m \rightarrow B_1, \dots, B_n \in \mathbf{GR}_3^-$ ,
- (2)  $A_1, \dots, A_m \rightarrow B_1, \dots, B_n \in \mathbf{GR}_2^-$ ,
- (3)  $A_1 \wedge \dots \wedge A_m \supset B_1 \vee \dots \vee B_n \in \mathbf{R}^-$ ,
- (4)  $A_1 \wedge \dots \wedge A_m \supset B_1 \vee \dots \vee B_n$  is valid in any Kripke model for  $\mathbf{R}^-$ .

**Corollary 3.17.** *If  $S \in \mathbf{GR}^-$ , then there exists a proof figure in  $\mathbf{GR}^-$  whose cuts are of the form of cuts occurring in page 5.*

Proof. Suppose that  $S \in \mathbf{GR}^-$ . By Lemma 3.1 and Corollary 3.16,  $S \in \mathbf{GR}_3^-$ . So, there exists a proof figure  $\mathcal{P}$  for  $S$  in  $\mathbf{GR}_3^-$ . Let  $\mathcal{Q}$  be the figure obtained from  $\mathcal{P}$  by replacing  $(\prec \rightarrow)$ ,  $(\text{tran})$  and  $(\text{lin})$  with the corresponding figure in page 5. We note that  $\mathcal{Q}$  is a proof figure for  $S$  in  $\mathbf{GR}^-$  and each cut is of the form of cuts occurring in page 5.  $\dashv$

## 4 A proof without extension lemma

In the previous section, we proved Theorem 3.15 using the extension lemma (Lemma 3.14) several times. So, for a sequent  $S_0 \notin \mathbf{GR}_3^-$ , a concrete Kripke model, in which the corresponding formula to  $S_0$  is not valid, is not clearly given. Here we show Theorem 3.15 without the extension lemma by modifying  $\mathcal{K}(S_0)$  in Definition 3.10. The proof also gives a concrete Kripke model, in which the corresponding formula to a sequent  $S_0 \notin \mathbf{GR}_3^-$  is not valid.

**Definition 4.1.** Let  $S_0$  be a sequent, which is not provable in  $\mathbf{GR}_3^-$ . We define a structure  $\mathcal{K}^*(S_0) = \langle \mathbf{W}(S_0), <, \models^* \rangle$  as follows:

- (1)  $<$  is as in Definition 3.10, we also write  $\alpha \leq \beta$  if  $\alpha < \beta$  or  $\alpha = \beta$ ,
- (2)  $\models^*$  is a valuation satisfying, in addition to the conditions in Definition 1.1(3),
  - (2.1)  $(\Gamma \rightarrow \Delta; \tau) \models^* p$  if and only if  $p \in \Gamma$ , for any propositional variable  $p$ ,
  - (2.2)  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box A \prec \Box B$  if and only if either one of the following three holds:
    - (2.2.1)  $\Box A \prec \Box B \in \Gamma$ ,
    - (2.2.2) there exists  $\alpha < (\Gamma \rightarrow \Delta; \tau)$  such that  $\alpha \models^* \Box A$  and  $\alpha \not\models^* \Box B$ ,
    - (2.2.3)  $\Box B \notin \text{Sub}(\Gamma \rightarrow \Delta)$  and there exist  $C$  and  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \leq (\Gamma \rightarrow \Delta; \tau)$  such that  $\Box A \prec \Box C \in \Gamma_1$ ,  $\Box A, \Box B, \Box C$  are true at  $(\Gamma_1 \rightarrow \Delta_1; \tau_1)$ , for any  $\alpha < (\Gamma_1 \rightarrow \Delta_1; \tau_1)$ ,  $\alpha \not\models^* \Box A$ ,  $\alpha \not\models^* \Box B$ ,  $\alpha \not\models^* \Box C$ ,
  - (2.3)  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box A \preceq \Box B$  if and only if either one of the following two holds:
    - (2.3.1)  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box A \prec \Box B$ ,
    - (2.3.2)  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box A$  and  $(\Gamma \rightarrow \Delta; \tau) \not\models^* \Box B \prec \Box A$ .

**Lemma 4.2.** *Let  $S_0$  be a sequent, which is not provable in  $\mathbf{GR}_3^-$ . Then for any  $A$  and for any  $(\Gamma \rightarrow \Delta; \tau) \in \mathbf{W}(S_0)$ ,*

- (1)  $A \in \Gamma$  *implies*  $(\Gamma \rightarrow \Delta; \tau) \models^* A$ ,
- (2)  $A \in \Delta$  *implies*  $(\Gamma \rightarrow \Delta; \tau) \not\models^* A$ .

*Proof.* We use an induction on  $A$ .

We only show the case that  $A$  is a  $\prec$ -formula or a  $\preceq$ -formula.

For the case that  $A = \Box B \prec \Box C$ . From Definition 4.1(2.2), (1) is clear. Suppose that

(3)  $\Box B \prec \Box C \in \Delta$ .

By Lemma 3.8(1),

(4)  $\Box B \prec \Box C \notin \Gamma$ .

By (3), we have

(5)  $\Box C \in \text{Sub}(\Gamma \rightarrow \Delta)$ .

So, we have only to show

(6) there does not exist  $\alpha \leq (\Gamma \rightarrow \Delta; \tau)$  such that  $\alpha \models^* \Box B$  and  $\alpha \not\models^* \Box C$ .

Suppose

(7) there exists  $(\Phi \rightarrow \Psi; \sigma) \leq (\Gamma \rightarrow \Delta; \tau)$  such that  $(\Phi \rightarrow \Psi; \sigma) \models^* \Box B$  and  $(\Phi \rightarrow \Psi; \sigma) \not\models^* \Box C$ .

Immediately, we have

(8)  $B \neq C$ .

Also by the induction hypothesis,

(9)  $\Box C \notin \Phi$  and  $\Box B \notin \Psi$ .

Since  $(\Phi \rightarrow \Psi; \sigma) \in \mathbf{W}(S_0)$ ,  $\Phi \rightarrow \Psi$  is saturated, and hence

(10)  $\Box C \preceq \Box B \notin \Phi$  and  $\Box C \prec \Box B \notin \Phi$ .

By Definition 3.3(11),

(11)  $\Box B \prec \Box C \in \Phi$ .

Using Lemma 3.9(2),

(12)  $\Box B \prec \Box C \in \Gamma$ .

This is contradictory to (3) and Lemma 3.8(1).

For the case that  $A = \Box B \preceq \Box C$ . Suppose

(13)  $\Box B \preceq \Box C \in \Gamma$ .

Since  $\Gamma \rightarrow \Delta$  is saturated,

(14)  $\Box B \in \Gamma$

and either

(15a)  $\Box B \prec \Box C \in \Gamma$  or (15b)  $\Box C \preceq \Box B \in \Gamma$ .

holds. If (15a) holds, then by the definition,

(16)  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box B \prec \Box C$ ,

and hence

(17)  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box B \preceq \Box C$ .

So, we assume that (15b) holds. By (15b), we have  $\Box C \preceq \Box B \notin \Delta$ , and hence

(18)  $\Box B \prec \Box C \notin \Gamma$ .

Similarly by (13), we have

(19)  $\Box C \prec \Box B \notin \Gamma$ .

Let it be that  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \leq (\Gamma \rightarrow \Delta; \tau)$ . By (18), (19) and Lemma 3.9(2),

(20)  $\Box C \prec \Box B \notin \Gamma_1$  and  $\Box B \prec \Box C \notin \Gamma_1$ .

By (13),

(21)  $\Box B, \Box C \in \text{Sub}(\Gamma \rightarrow \Delta) \subseteq \text{Sub}(\Gamma_1 \rightarrow \Delta_1)$ .

Since  $\Gamma_1 \rightarrow \Delta_1$  is saturated,

(22) either  $\Box C \preceq \Box B, \Box B \preceq \Box C, \Box B, \Box C \in \Gamma_1$  or  $\Box B, \Box C \in \Delta_1$  holds.

By the induction hypothesis,

(23) at  $(\Gamma_1 \rightarrow \Delta_1; \tau_1)$ , both of  $\Box B$  and  $\Box C$  are true or none of  $\Box B$  and  $\Box C$  is true.

Hence

(24) for any  $\alpha \leq (\Gamma \rightarrow \Delta; \tau)$ ,  $\alpha \not\models^* \Box C$  or  $\alpha \models^* \Box B$ .

Also by (13), we have

(25)  $\Box B \in \text{Sub}(\Gamma \rightarrow \Delta)$ .  
 By (18), (24) and (25), we have  
 (26)  $(\Gamma \rightarrow \Delta; \tau) \not\models^* \Box C \prec \Box B$ .  
 Using (14), we obtain  
 (26)  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box B \preceq \Box C$ .

Suppose  
 (27)  $\Box B \preceq \Box C \in \Delta$ .  
 Then we have  
 (28)  $\Box B, \Box C \in \text{Sub}(\Gamma \rightarrow \Delta)$ .  
 and  
 (29)  $\Box B \preceq \Box C \notin \Gamma$ .  
 Since  $\Gamma \rightarrow \Delta$  is saturated,  
 (30)  $\Box B \prec \Box C \notin \Gamma$ ,  
 and

(31) either  $\Box B, \Box C \in \Delta$  or  $\Box C \prec \Box B \in \Gamma$  holds.  
 Let it be that  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \leq (\Gamma \rightarrow \Delta; \tau)$ . Then by (29), (30) and Lemma 3.9(2), we have  
 (32)  $\Box B \preceq \Box C \notin \Gamma_1$  and  $\Box B \prec \Box C \notin \Gamma_1$ .

By (28), we have  
 (33)  $\Box B, \Box C \in \text{Sub}(\Gamma \rightarrow \Delta) \subseteq \text{Sub}(\Gamma_1 \rightarrow \Delta_1)$ .  
 Since  $\Gamma_1 \rightarrow \Delta_1$  is saturated,  
 (33) either  $\Box B, \Box C \in \Delta_1$  or  $\Box C \prec \Box B, \Box C \in \Gamma_1$  holds.  
 By the induction hypothesis,  
 (33) either  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \not\models^* \Box B$  or  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \models^* \Box C$

Hence  
 (34) for any  $\alpha \leq (\Gamma \rightarrow \Delta; \tau)$ ,  $\alpha \not\models^* \Box B$  or  $\alpha \models^* \Box C$ .  
 Using (28) and (30), we have

(35)  $(\Gamma \rightarrow \Delta; \tau) \not\models^* \Box B \prec \Box C$ .  
 On the other hand, by (31), the induction hypothesis and a result in the above case,  
 (36) either  $(\Gamma \rightarrow \Delta; \tau) \not\models^* \Box B$  or  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box C \preceq \Box B$  holds.

Hence we obtain  
 (37)  $(\Gamma \rightarrow \Delta; \tau) \not\models^* \Box B \preceq \Box C$ . ⊥

**Corollary 4.3.** *Let  $A_1, \dots, A_m \rightarrow B_1, \dots, B_n$  be a sequent, which is not provable in  $\mathbf{GR}_3^-$ . Then in  $\mathcal{K}^*(A_1, \dots, A_m \rightarrow B_1, \dots, B_n)$ ,*

$$(\text{sat}(A_1, \dots, A_m \rightarrow B_1, \dots, B_n); [\ ] ) \not\models^* A_1 \wedge \dots \wedge A_m \supset B_1 \vee \dots \vee B_n.$$

**Lemma 4.4.** *Let  $S_0$  be a sequent, which is not provable in  $\mathbf{GR}_3^-$ . Then  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box A \preceq \Box B$  implies  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box A$ .*

Proof. Suppose that  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box A \preceq \Box B$ . If  $(\Gamma \rightarrow \Delta; \tau) \not\models^* \Box A \prec \Box B$ , then by the definition we have  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box A$ . So, we assume that  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box A \prec \Box B$ . Then either one of the following three holds:

- (1)  $\Box A \prec \Box B \in \Gamma$ ,
- (2) there exists  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \leq (\Gamma \rightarrow \Delta; \tau)$  such that  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \models^* \Box A$  and  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \not\models^* \Box B$ ,
- (3) there exist a formula  $C$  and  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \leq (\Gamma \rightarrow \Delta; \tau)$  such that
  - $\Box A \prec \Box C \in \Gamma_1$ ,
  - $\Box B \notin \text{Sub}(\Gamma_1 \rightarrow \Delta_1)$ ,
  - $\Box A, \Box B, \Box C$  are true at  $(\Gamma_1 \rightarrow \Delta_1; \tau_1)$ ,
  - for any  $\alpha < (\Gamma_1 \rightarrow \Delta_1; \tau_1)$ ,  $\alpha \not\models^* \Box A$ ,  $\alpha \not\models^* \Box B$ ,  $\alpha \not\models^* \Box C$ .

If (1) holds, then  $\Box A, \Box A \preceq \Box B \in \Gamma$ . Using Lemma 4.2, we have  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box A$ .



If either (2) or (3) holds, then there exists  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \leq (\Gamma \rightarrow \Delta; \tau)$  such that  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \models^* \Box A$ . So,  $(\Gamma_2 \rightarrow \Delta_2; \tau_2) \models^* A$  for any  $(\Gamma_2 \rightarrow \Delta_2; \tau_2) \in (\Gamma_1 \rightarrow \Delta_1; \tau_1) \uparrow$ . So,  $(\Gamma_2 \rightarrow \Delta_2; \tau_2) \models^* A$  for any  $(\Gamma_2 \rightarrow \Delta_2; \tau_2) \in (\Gamma \rightarrow \Delta; \tau) \uparrow$ . Hence  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box A$ .  $\dashv$

**Lemma 4.5.** *Let  $S_0$  be a sequent, which is not provable in  $\mathbf{GR}_3^-$ . Then  $\mathcal{K}^*(S_0)$  is a Kripke pseudo-model for  $\mathbf{R}^-$ .*

Proof. By Lemma 3.13.  $\dashv$

**Lemma 4.6.** *Let  $S_0$  be a sequent, which is not provable in  $\mathbf{GR}_3^-$ . Then for any  $D \in \Sigma$ ,  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \models^* D$  and  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) < (\Gamma_2 \rightarrow \Delta_2; \tau_2)$  imply  $(\Gamma_2 \rightarrow \Delta_2; \tau_2) \models^* D$ .*

Proof. Suppose that  $D$  is either a  $\preceq$ -formula or a  $\prec$ -formula,  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \models^* D$  and  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) < (\Gamma_2 \rightarrow \Delta_2; \tau_2)$ .

If  $D$  is  $\Box$ -formula, then the lemma is clear by the definition of  $\models^*$ .

If  $D = \Box A \prec \Box B$ , then either one of the following three holds:

- (1)  $\Box A \prec \Box B \in \Gamma_1$ ,
- (2) there exists  $\alpha \leq (\Gamma_1 \rightarrow \Delta_1; \tau_1)$  such that  $\alpha \models^* \Box A$  and  $\alpha \not\models^* \Box B$ ,
- (3)  $\Box B \notin \text{Sub}(\Gamma_1 \rightarrow \Delta_1)$  and there exist  $C$  and  $(\Gamma_3 \rightarrow \Delta_3; \tau_3) \leq (\Gamma_1 \rightarrow \Delta_1; \tau_1)$  such that
  - $\Box A \prec \Box C \in \Gamma_3$ ,
  - $\Box A, \Box B, \Box C$  are true at  $(\Gamma_3 \rightarrow \Delta_3; \tau_3)$ ,
  - for any  $\alpha < (\Gamma_3 \rightarrow \Delta_3; \tau_3)$ ,  $\alpha \not\models^* \Box A$ ,  $\alpha \not\models^* \Box B$ ,  $\alpha \not\models^* \Box C$ .

If (1) holds, then by Lemma 3.9(2),  $\Box A \prec \Box B \in \Gamma_2$ , and by Lemma 4.2,  $(\Gamma_2 \rightarrow \Delta_2; \tau_2) \models^* D$ . If (2) holds, then we have  $\alpha < (\Gamma_2 \rightarrow \Delta_2; \tau_2)$ , and hence  $(\Gamma_2 \rightarrow \Delta_2; \tau_2) \models^* D$ . If (3) holds, then similarly,  $(\Gamma_3 \rightarrow \Delta_3; \tau_3) < (\Gamma_2 \rightarrow \Delta_2; \tau_2)$ . We also note that  $\Box B \notin \text{Sub}(\Gamma_1 \rightarrow \Delta_1) \supseteq \text{Sub}(\Gamma_2 \rightarrow \Delta_2)$ . Hence we obtain  $(\Gamma_2 \rightarrow \Delta_2; \tau_2) \models^* D$ .

If  $D = \Box A \preceq \Box B$ , then either one of the following two holds:

- (4)  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \models^* \Box A \prec \Box B$ ,
- (5)  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \models^* \Box A$  and  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \not\models^* \Box B \prec \Box A$ .

If (4) holds, then by the case above,  $(\Gamma_2 \rightarrow \Delta_2; \tau_2) \models^* \Box A \prec \Box B$ , and hence  $(\Gamma_2 \rightarrow \Delta_2; \tau_2) \models^* \Box A \preceq \Box B$ . Assume that (5) holds. From the definition, we have  $(\Gamma_2 \rightarrow \Delta_2; \tau_2) \models^* \Box A$ . So, it is sufficient to show  $(\Gamma_2 \rightarrow \Delta_2; \tau_2) \not\models^* \Box B \prec \Box A$ . Suppose that  $(\Gamma_2 \rightarrow \Delta_2; \tau_2) \models^* \Box B \prec \Box A$ . Then either one of the following three holds:

- (6)  $\Box B \prec \Box A \in \Gamma_2$  and  $A \neq B$ ,
- (7) there exists  $(\Gamma_3 \rightarrow \Delta_3; \tau_3) \leq (\Gamma_2 \rightarrow \Delta_2; \tau_2)$  such that  $(\Gamma_3 \rightarrow \Delta_3; \tau_3) \models^* \Box B$  and  $(\Gamma_3 \rightarrow \Delta_3; \tau_3) \not\models^* \Box A$ ,
- (8)  $\Box B \notin \text{Sub}(\Gamma_2 \rightarrow \Delta_2)$  and there exist a formula  $C$  and  $(\Gamma_3 \rightarrow \Delta_3; \tau_3) \leq (\Gamma_2 \rightarrow \Delta_2; \tau_2)$  such that
  - $\Box A \prec \Box C \in \Gamma_3$ ,
  - $\Box A, \Box B, \Box C$  are true at  $(\Gamma_3 \rightarrow \Delta_3; \tau_3)$ ,
  - for any  $\alpha < (\Gamma_3 \rightarrow \Delta_3; \tau_3)$ ,  $\alpha \not\models^* \Box A$ ,  $\alpha \not\models^* \Box B$ ,  $\alpha \not\models^* \Box C$ .

If (6) holds, then  $\Box A \preceq \Box B \in \Delta_2$  and  $\Box A \preceq \Box B \in \text{Sub}^+(\Gamma_2 \rightarrow \Delta_2) \subseteq \text{Sub}^+(\Gamma_1 \rightarrow \Delta_1)$ . On the other hand, by (5) and Lemma 4.2, we have  $\Box A \notin \Delta_1$  and  $\Box B \prec \Box A \notin \Gamma_1$ . Hence  $\Box A \preceq \Box B \in \Gamma_1$ . Using Lemma 3.9(2),  $\Box A \preceq \Box B \in \Gamma_2$ . This is in contradiction with  $\Box A \preceq \Box B \in \Delta_2$  and Lemma 3.8.

If (7) holds, then by (5),  $(\Gamma_3 \rightarrow \Delta_3; \tau_3) < (\Gamma_1 \rightarrow \Delta_1; \tau_1)$ . So, we have  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \models^* \Box B \prec \Box A$ . This is in contradiction with (5).

If (8) holds, then by (5),  $(\Gamma_3 \rightarrow \Delta_3; \tau_3) < (\Gamma_1 \rightarrow \Delta_1; \tau_1)$ . So, if  $\Box B \notin \text{Sub}(\Gamma_1 \rightarrow \Delta_1)$ , then we have  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \models^* \Box B \prec \Box A$ , and this is in contradiction with (5). Assume that  $\Box B \in \text{Sub}(\Gamma_1 \rightarrow \Delta_1)$ . By (8),  $\Box A \in \text{Sub}(\Gamma_3 \rightarrow \Delta_3) \subseteq \text{Sub}(\Gamma_1 \rightarrow \Delta_1)$ . By (5), Lemma 4.2 and Lemma 3.9(3), we have  $\Box A \preceq \Box B \in \Gamma_1 \cap \Sigma \subseteq \Gamma_2$ , similarly to the case that (6) holds. Hence  $\Box B \in \text{Sub}(\Gamma_2 \rightarrow \Delta_2)$ . This is contradictory to (8).  $\dashv$

**Lemma 4.7.** *Let  $S_0$  be a sequent, which is not provable in  $\mathbf{GR}_3^-$  and let be that  $\alpha \models^* \Box A \preceq \Box B$ . Then for any  $\beta \leq \alpha$ ,  $\beta \models^* \Box B$  implies  $\beta \models^* \Box A$ .*

Proof. We put  $\alpha = (\Gamma \rightarrow \Delta; \tau)$ . Then either one of the following four holds:

- (1)  $\alpha \models^* \Box A$  and  $\alpha \not\models^* \Box B \prec \Box A$ .
- (2)  $\Box A \prec \Box B \in \Gamma$ ,
- (3) there exists  $\gamma \leq \alpha$  such that  $\gamma \models^* \Box A$  and  $\gamma \not\models^* \Box B$ ,
- (4)  $\Box B \notin \text{Sub}(\Gamma \rightarrow \Delta)$  and there exist  $C$  and  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \leq (\Gamma \rightarrow \Delta; \tau)$  such that
  - $\Box A \prec \Box C \in \Gamma_1$ ,
  - $\Box A, \Box B, \Box C$  are true at  $(\Gamma_1 \rightarrow \Delta_1; \tau_1)$ ,
  - for any  $\delta < (\Gamma_1 \rightarrow \Delta_1; \tau_1)$ ,  $\delta \not\models^* \Box A$ ,  $\delta \not\models^* \Box B$ ,  $\delta \not\models^* \Box C$ .

If (1) holds, then we obtain the lemma by the definition.

If (2) holds, then  $\Box B \preceq \Box A \in \Delta$ . Since  $\Gamma \rightarrow \Delta \notin \mathbf{GR}_3^-$ ,  $\Box B \preceq \Box A \notin \Gamma$ . Hence  $\Box B \prec \Box A \notin \Gamma$ . Using Lemma 3.9(2) for any  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \leq (\Gamma \rightarrow \Delta; \tau)$  we have  $\Box B \prec \Box A \notin \Gamma_1$ . Using  $\Box A, \Box B \in \text{Sub}(\Gamma \rightarrow \Delta) \subseteq \text{Sub}(\Gamma_1 \rightarrow \Delta_1)$ , we have either  $\Box B \in \Delta_1$ ,  $\Box A \prec \Box B \in \Gamma_1$  or  $\Box A \preceq \Box B \in \Gamma_1$ . Hence  $\Box B \in \Delta_1$  or  $\Box A \in \Gamma_1$ . Using Lemma 4.2,  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \not\models^* \Box B$  or  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \models^* \Box A$ .

If (3) holds, then by Lemma 4.6,  $\gamma_1 \models^* \Box A$  for any  $\gamma_1 \in \gamma \uparrow$  and  $\gamma_2 \not\models^* \Box B$  for any  $\gamma_2 \leq \gamma$ . We note that for any  $\beta \leq \alpha$ , either  $\beta \leq \gamma$  or  $\gamma < \beta$ . So, we obtain the lemma.

If (4) holds, then  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \leq \alpha$ . Suppose that  $\beta \leq \alpha$ . By Lemma 4.5,  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \leq \beta$  or  $\beta < (\Gamma_1 \rightarrow \Delta_1; \tau_1)$ . By (4),  $\beta \models^* \Box A$  if  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \leq \beta$ ; and  $\beta \not\models^* \Box B$  if  $\beta < (\Gamma_1 \rightarrow \Delta_1; \tau_1)$ . So,  $\beta \models^* \Box B$  implies  $\beta \models^* \Box A$ .  $\dashv$

**Lemma 4.8.** *Let  $S_0$  be a sequent, which is not provable in  $\mathbf{GR}_3^-$ . then  $\mathcal{K}^*(S_0)$  is a Kripke model for  $\mathbf{R}^-$ .*

Proof. By Lemma 4.5 and Lemma 4.6, it is sufficient to show the condition (2) in Definition 1.2.

We divide the cases.

The case that  $D$  is A5 (i.e.,  $D = (\Box A \preceq \Box B) \supset \Box A$ ) is shown by Lemma 4.4.

For the case that  $D$  is A6 (i.e.,  $D = ((\Box A \preceq \Box B) \wedge (\Box B \preceq \Box C)) \supset (\Box A \preceq \Box C)$ ). If either  $A = B$  or  $B = C$ , then  $(\Gamma \rightarrow \Delta; \tau) \models D$  is clear. So, we assume that  $A \neq B$  and  $B \neq C$ . Suppose that

- (1)  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box A \preceq \Box B$ ,
- (2)  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box B \preceq \Box C$  and
- (3)  $(\Gamma \rightarrow \Delta; \tau) \not\models^* \Box A \preceq \Box C$ .

By (3), we have

- (4)  $(\Gamma \rightarrow \Delta; \tau) \not\models^* \Box A \prec \Box C$  and
- (5)  $(\Gamma \rightarrow \Delta; \tau) \not\models^* \Box A$  or  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box C \prec \Box A$ .

By (1) and Lemma 4.4, we have

- (6)  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box A$ .

Using (5), we have  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box C \prec \Box A$ . So, either one of the following three holds:

- (3a)  $\Box C \prec \Box A \in \Gamma$ ,
- (3b) there exists  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \leq (\Gamma \rightarrow \Delta; \tau)$  such that  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \models^* \Box C$  and  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \not\models^* \Box A$ ,

- (3c)  $\Box A \notin \text{Sub}(\Gamma \rightarrow \Delta)$ , and there exist a formula  $D$  and  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \leq (\Gamma \rightarrow \Delta; \tau)$  such that
  - $\Box C \prec \Box D \in \Gamma_1$ ,
  - $\Box A, \Box C, \Box D$  are true at  $(\Gamma_1 \rightarrow \Delta_1; \tau_1)$ ,
  - for any  $\alpha < (\Gamma_1 \rightarrow \Delta_1; \tau_1)$ ,  $\alpha \not\models^* \Box A$ ,  $\alpha \not\models^* \Box C$ ,  $\alpha \not\models^* \Box D$ .

By (1), either one of the following four holds:

- (1a)  $(\Gamma \rightarrow \Delta; \tau) \not\models^* \Box B \prec \Box A$ .
- (1b)  $\Box A \prec \Box B \in \Gamma$ ,
- (1c) there exists  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \leq (\Gamma \rightarrow \Delta; \tau)$  such that  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \models^* \Box A$  and  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \not\models^* \Box B$ ,

- (1d)  $\Box B \notin \text{Sub}(\Gamma \rightarrow \Delta)$ , and there exist a formula  $D$  and  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \leq (\Gamma \rightarrow \Delta; \tau)$  such that
  - $\Box A \prec \Box D \in \Gamma_1$ ,
  - $\Box A, \Box B, \Box D$  are true at  $(\Gamma_1 \rightarrow \Delta_1; \tau_1)$ ,
  - for any  $\alpha < (\Gamma_1 \rightarrow \Delta_1; \tau_1)$ ,  $\alpha \not\models^* \Box A$ ,  $\alpha \not\models^* \Box B$ ,  $\alpha \not\models^* \Box D$ .

By (2), either one of the following four holds:

- (2a)  $(\Gamma \rightarrow \Delta; \tau) \not\models^* \Box C \prec \Box B$ .  
(2b)  $\Box B \prec \Box C \in \Gamma$ ,  
(2c) there exists  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \leq (\Gamma \rightarrow \Delta; \tau)$  such that  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \models^* \Box B$  and  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \not\models^* \Box C$ ,  
(2d)  $\Box C \notin \text{Sub}(\Gamma \rightarrow \Delta)$  and there exist a formula  $D$  and  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \leq (\Gamma \rightarrow \Delta; \tau)$  such that  
 $\Box B \prec \Box D \in \Gamma_1$ ,  
 $\Box C, \Box C, \Box D$  are true at  $(\Gamma_1 \rightarrow \Delta_1; \tau_1)$ ,  
for any  $\alpha < (\Gamma_1 \rightarrow \Delta_1; \tau_1)$ ,  $\alpha \not\models^* \Box B$ ,  $\alpha \not\models^* \Box C$ ,  $\alpha \not\models^* \Box D$ .

We divide the subcases.

The subcase that (3a) and  $\Box B \in \text{Sub}(\Gamma \rightarrow \Delta)$  hold. We note that  $\Box A, \Box B, \Box C \in \text{Sub}(\Gamma \rightarrow \Delta)$ . By (1) and (2) and Lemma 4.2, we have  $\Box A \preceq \Box B \notin \Delta$  and  $\Box B \preceq \Box C \notin \Delta$ . Since  $\Gamma \rightarrow \Delta$  is saturated, we have  $\Box A \preceq \Box C \in \Gamma$ . Using Lemma 4.2,  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box A \preceq \Box C$ . This is contradictory to (3).

The subcase that (3a), (2a) and  $\Box B \notin \text{Sub}(\Gamma \rightarrow \Delta)$  hold. By (3a), we have  $\Box A \preceq \Box C \in \Delta$ , and hence  $\Box A \preceq \Box C \notin \Gamma$  and  $\Box A \prec \Box C \notin \Gamma$ . Let  $(\Gamma_1 \rightarrow \Delta_1; \tau_1)$  be a world in  $(\Gamma \rightarrow \Delta; \tau) \downarrow \cup \{(\Gamma \rightarrow \Delta; \tau)\}$ . By Lemma 3.9(2),  $\Box A \preceq \Box C \notin \Gamma_1$  and  $\Box A \prec \Box C \notin \Gamma_1$ . So, either  $\Box C \prec \Box A, \Box C \preceq \Box A, \Box C \in \Gamma_1$  or  $\Box A, \Box C \in \Delta_1$ . Using Lemma 4.2, either  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \models^* \Box C$  or  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \not\models^* \Box A$ . Using (1), (2) and Lemma 4.7, either each of  $\Box A, \Box B, \Box C$  is true or none of  $\Box A, \Box B, \Box C$  is true at  $(\Gamma_1 \rightarrow \Delta_1; \tau_1)$ . By Lemma 4.5,  $((\Gamma \rightarrow \Delta; \tau) \downarrow \cup \{(\Gamma \rightarrow \Delta; \tau)\}, <)$  is a linear ordered set and is not empty. So, there exists the minimal element  $(\Gamma_2 \rightarrow \Delta_2; \tau_2)$  that makes each of  $\Box A, \Box B, \Box C$  true. We note that

- $(\Gamma_2 \rightarrow \Delta_2; \tau_2) \leq (\Gamma \rightarrow \Delta; \tau)$   
 $\Box C \prec \Box A \in \Gamma_2$ ,  
 $\Box C, \Box B, \Box A$  are true at  $(\Gamma_2 \rightarrow \Delta_2; \tau_2)$ ,  
for any  $\alpha < (\Gamma_2 \rightarrow \Delta_2; \tau_2)$ ,  $\alpha \not\models^* \Box C$ ,  $\alpha \not\models^* \Box B$ ,  $\alpha \not\models^* \Box A$ .

Also by the assumption,  $\Box B \notin \text{Sub}(\Gamma \rightarrow \Delta)$ . So, we have  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box C \prec \Box B$ . This is contradictory to (2a).

The subcase that (3a) and (2b) hold. We note that  $\Box A, \Box B, \Box C \in \text{Sub}(\Gamma \rightarrow \Delta)$ . So, this subcase resolves into the first subcase.

The subcase that (3a) and (2c) hold. By (2c), we have  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \models^* \Box B$  and  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \not\models^* \Box C$ . Using (1) and Lemma 4.7,  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \models^* \Box A$ . Using Lemma 4.2,  $\Box A \notin \Delta_1$  and  $\Box C \notin \Gamma_1$ . So,  $\Box C \prec \Box A \notin \Gamma_1$  and  $\Box C \preceq \Box A \notin \Gamma_1$ . By (3a), we have  $\Box A, \Box C \in \text{Sub}(\Gamma \rightarrow \Delta) \subseteq \text{Sub}(\Gamma_1 \rightarrow \Delta_1)$ . Hence  $\Box A \prec \Box C \in \Gamma_1$ . Using Lemma 3.9(2),  $\Box A \prec \Box C \in \Gamma$ . This is contradictory to (3a).

The subcase that (3a) and (2d) hold. By (3a), we have  $\Box C \prec \Box D \in \Gamma_1$ . Using Lemma 3.9(2),  $\Box C \prec \Box D \in \Gamma$ . So,  $\Box C \in \text{Sub}(\Gamma \rightarrow \Delta)$ . This is contradictory to (2d).

The subcase that (3b) holds. Using (1) and Lemma 4.7,  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \not\models^* \Box B$ . Using (2) and Lemma 4.7,  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \not\models^* \Box A$ . This is contradictory to (3b).

The subcase that (3c), (1a) and  $\Box B \in \text{Sub}(\Gamma \rightarrow \Delta)$  hold. We note that  $\Box B \in \text{Sub}(\Gamma \rightarrow \Delta) \subseteq \text{Sub}(\Gamma_1 \rightarrow \Delta_1)$  hold. Using (3c), we have  $\Box B, \Box C, \Box D \in \text{Sub}(\Gamma_1 \rightarrow \Delta_1)$ . On the other hand, by (2) and Lemma 4.2,  $\Box B \preceq \Box C \notin \Delta$ , and hence  $\Box C \prec \Box B \notin \Gamma$ . Using Lemma 4.6,  $\Box C \prec \Box B \notin \Gamma_1$ . Also by (3c) and Lemma 4.2,  $\Box C \notin \Delta_1$ . Hence we have either  $\Box B \prec \Box C, \Box B \preceq \Box C \in \Gamma_1$  or  $\Box B \preceq \Box C, \Box C \preceq \Box B \in \Gamma_1$ . So, we have  $\Box B \preceq \Box C \in \Gamma_1$ . By (3c),  $\Box C \prec \Box D, \Box C \preceq \Box D \in \Gamma_1$ . So,  $\Box B \preceq \Box D \in \Gamma_1$ . Also,  $\Box B \in \Gamma_1$  and that  $\Box D \prec \Box B \in \Gamma_1$  implies  $\Box B \preceq \Box D \in \Delta_1$ . Hence  $\Box B \notin \Delta_1$  and  $\Box D \prec \Box B \notin \Gamma_1$ . Since  $\Gamma_1 \rightarrow \Delta_1$  is saturated, we have either  $\Box B \prec \Box D \in \Gamma_1$  or  $\Box D \preceq \Box B \in \Gamma_1$ . If  $\Box D \preceq \Box B \in \Gamma_1$ , then by  $\Box B \preceq \Box C \in \Gamma_1$ , we have  $\Box D \preceq \Box C \in \Gamma_1$ , and hence  $\Box C \prec \Box D \in \Delta_1$ . This is in contradiction with  $\Box C \prec \Box D \in \Gamma_1$ . Assume that  $\Box B \prec \Box D \in \Gamma_1$ . By  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \models^* \Box C$ , (2) and Lemma 4.7, we have  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \models^* \Box B$ . By  $\alpha \not\models^* \Box A$  for any  $\alpha < (\Gamma_1 \rightarrow \Delta_1; \tau_1)$ , (1) and Lemma 4.7, we have  $\alpha \not\models^* \Box B$  for any  $\alpha < (\Gamma_1 \rightarrow \Delta_1; \tau_1)$ . So, using (3c), we have  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box B \prec \Box A$ . This is in contradiction with (1a).

The subcase that (3c)[, (1a)],  $\Box B \notin \text{Sub}(\Gamma \rightarrow \Delta)$  and (2a) hold. By (3c), (1), (2) and Lemma 4.7, we have  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \models^* \Box B$ .  $\alpha \not\models \Box B$  for any  $\alpha < (\Gamma_1 \rightarrow \Delta_1; \tau_1)$ . Using (3c) and  $\Box B \notin \text{Sub}(\Gamma \rightarrow \Delta)$ , we have  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box C \prec \Box B$ . This is in contradiction with (2a).

The subcase that  $[(3c), (1a)]$ ,  $\Box B \notin \text{Sub}(\Gamma \rightarrow \Delta)$  and (2b) hold. By (2b),  $\Box B \in \text{Sub}(\Gamma \rightarrow \Delta)$ , which is in contradiction with  $\Box B \notin \text{Sub}(\Gamma \rightarrow \Delta)$ .

The subcase that  $(3c)[(1a), \Box B \notin \text{Sub}(\Gamma \rightarrow \Delta)]$  and (2c) hold. By (2c), there exists  $\alpha \leq (\Gamma \rightarrow \Delta; \tau)$  such that  $\alpha \models^* \Box B$  and  $\alpha \not\models^* \Box C$ . By (3c), there exists  $\beta \leq (\Gamma \rightarrow \Delta; \tau)$  such that  $\beta \models^* \Box C$  and  $\gamma \not\models^* \Box A$  for any  $\gamma \in \beta \downarrow$ . Using Lemma 4.5,  $\alpha < \beta$  or  $\beta \leq \alpha$ . By  $\alpha \not\models^* \Box C$ ,  $\beta \models^* \Box C$  and Lemma 4.6, we have  $\alpha < \beta$ . Since  $\gamma \not\models^* \Box A$  for any  $\gamma \in \beta \downarrow$ ,  $\alpha \not\models^* \Box A$ . Using (1) and Lemma 4.7, we have  $\alpha \not\models^* \Box B$ . This is in contradiction in  $\alpha \models^* \Box B$ .

The subcase that  $(3c)[(1a), \Box B \notin \text{Sub}(\Gamma \rightarrow \Delta)]$  and (2d) hold. By (3c) and Lemma 3.9(2), we have  $\Box A \prec \Box C \in \Gamma_1 \subseteq \Gamma$ , and hence  $\Box A \in \text{Sub}(\Gamma \rightarrow \Delta)$ . This is in contradiction with  $\Box C \notin \text{Sub}(\Gamma \rightarrow \Delta)$  from (2d).

The subcase that (3c) and (1b) hold. By (3c), we have  $\Box A \notin \text{Sub}(\Gamma \rightarrow \Delta)$ . By (1b), we have  $\Box A \in \text{Sub}(\Gamma \rightarrow \Delta)$ . This is a contradiction.

The subcase that (3c) and (1c) hold. By (1c), there exists  $\alpha \leq (\Gamma \rightarrow \Delta; \tau)$  such that  $\alpha \models^* \Box A$  and  $\alpha \not\models^* \Box B$ . By (3c), there exists  $\beta \leq (\Gamma \rightarrow \Delta; \tau)$  such that  $\beta \models^* \Box C$  and  $\gamma \not\models^* \Box A$  for any  $\gamma \in \beta \downarrow$ . Using Lemma 4.5,  $\alpha < \beta$  or  $\beta \leq \alpha$ . Since  $\alpha \models^* \Box A$  and  $\gamma \not\models^* \Box A$  for any  $\gamma \in \beta \downarrow$ , we have  $\beta \leq \alpha$ . Using Lemma 4.6, we have  $\alpha \models^* \Box C$ . Using (2) and Lemma 4.7, we have  $\alpha \models^* \Box B$ . This is in contradiction in  $\alpha \not\models^* \Box B$ .

The subcase that (3c) and (1d) hold. By (1d) and Lemma 3.9(2), we have  $\Box A \prec \Box D \in \Gamma_1 \subseteq \Gamma$ , and hence  $\Box A \in \text{Sub}(\Gamma \rightarrow \Delta)$ . This is in contradiction with  $\Box A \notin \text{Sub}(\Gamma \rightarrow \Delta)$  from (3c).

For the case that  $D$  is A7 (i.e.,  $D = (\Box A \vee \Box B) \supset ((\Box A \preceq \Box B) \vee (\Box B \prec \Box A))$ ). Suppose that  $(\Gamma \rightarrow \Delta; \tau) \models \Box A \vee \Box B$ , i. e., either  $\Box A$  or  $\Box A$  is true at the pair.

If  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box A$ , then from the definition,  $(\Gamma \rightarrow \Delta; \tau) \not\models^* \Box B \prec \Box A$  implies  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box A \preceq \Box B$ . So,  $(\Gamma \rightarrow \Delta; \tau) \models^* (\Box A \preceq \Box B) \vee (\Box B \prec \Box A)$ .

If  $(\Gamma \rightarrow \Delta; \tau) \not\models^* \Box A$ , then  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box B$ . So, there exists  $\alpha \leq (\Gamma \rightarrow \Delta; \tau)$  such that  $\alpha \models^* \Box B$ .  $\alpha \not\models^* \Box A$ , and hence  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box B \prec \Box A$ .

For the case that  $D$  is A8 (i.e.,  $D = (\Box A \prec \Box B) \supset (\Box A \preceq \Box B)$ ). From the definition,  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box A \prec \Box B$  implies  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box A \preceq \Box B$ .

For the case that  $D$  is A9 (i.e.,  $D = ((\Box A \preceq \Box B) \wedge (\Box B \prec \Box A)) \supset \perp$ ). Suppose that

(7)  $(\Gamma \rightarrow \Delta; \tau) \models \Box A \preceq \Box B$  and

(8)  $(\Gamma \rightarrow \Delta; \tau) \models \Box B \prec \Box A$ .

By (8), either one of the following three holds:

(8a)  $\Box B \prec \Box A \in \Gamma$ ,

(8b) there exists  $\alpha \leq (\Gamma \rightarrow \Delta; \tau)$  such that  $\alpha \models^* \Box B$  and  $\alpha \not\models^* \Box A$ ,

(8c)  $\Box A \notin \text{Sub}(\Gamma \rightarrow \Delta)$  and there exist  $C$  and  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \leq (\Gamma \rightarrow \Delta; \tau)$  such that

$\Box B \prec \Box C \in \Gamma_1$ ,

$\Box A, \Box B, \Box C$  are true at  $(\Gamma_1 \rightarrow \Delta_1; \tau_1)$ ,

for any  $\alpha < (\Gamma_1 \rightarrow \Delta_1; \tau_1)$ ,  $\alpha \not\models^* \Box A$ ,  $\alpha \not\models^* \Box B$ ,  $\alpha \not\models^* \Box C$ .

If (8a) holds, then  $\Box A \preceq \Box B \in \Delta$  since  $\Gamma \rightarrow \Delta$  is saturated. Using Lemma 4.2,  $(\Gamma \rightarrow \Delta; \tau) \not\models^* \Box A \preceq \Box B$ . This is in contradiction with (7).

By (7) and Lemma 4.7, for any  $\alpha \leq (\Gamma \rightarrow \Delta; \tau)$ ,  $\alpha \models^* \Box B$  implies  $\alpha \not\models^* \Box A$ . This is in contradiction with (8b). So, (8b) does not hold.

Assume that (8c) holds. By (7), either one of the following four holds:

(7a)  $(\Gamma \rightarrow \Delta; \tau) \models^* \Box A$  and  $(\Gamma \rightarrow \Delta; \tau) \not\models^* \Box B \prec \Box A$ ,

(7b)  $\Box A \prec \Box B \in \Gamma$ ,

(7c) there exists  $\alpha \leq (\Gamma \rightarrow \Delta; \tau)$  such that  $\alpha \models^* \Box A$  and  $\alpha \not\models^* \Box B$ ,

(7d)  $\Box B \notin \text{Sub}(\Gamma \rightarrow \Delta)$  and there exist  $C$  and  $(\Gamma_1 \rightarrow \Delta_1; \tau_1) \leq (\Gamma \rightarrow \Delta; \tau)$  such that

$\Box A \prec \Box C \in \Gamma_1$ ,

$\Box A, \Box B, \Box C$  are true at  $(\Gamma_1 \rightarrow \Delta_1; \tau_1)$ ,

for any  $\alpha < (\Gamma_1 \rightarrow \Delta_1; \tau_1)$ ,  $\alpha \not\models^* \Box A$ ,  $\alpha \not\models^* \Box B$ ,  $\alpha \not\models^* \Box C$ .

(7a) is in contradiction with (8). (7b) is in contradiction with  $\Box A \notin \text{Sub}(\Gamma \rightarrow \Delta)$  from (8c). By (8) and

Lemma 4.7, for any  $\alpha \leq (\Gamma \rightarrow \Delta; \tau)$ ,  $\alpha \models^* \Box A$  implies  $\alpha \not\models^* \Box B$ , which is in contradiction with (7c). If (7d) holds, then  $\Box A \prec \Box C \in \Gamma_1$ , using Lemma 3.9(2),  $\Box A \prec \Box C \in \Gamma \subseteq \text{Sub}(\Gamma \rightarrow \Delta)$ , which is in contradiction with  $\Box A \notin \text{Sub}(\Gamma \rightarrow \Delta)$  from (8c).

From Corollary 4.3 and Lemma 4.8, we obtain Theorem 3.15.

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